

The Design of an IBM
MBC Interpreter

BYTE NYBBLES

The Design of an M6800 LISP Interpreter

S Tucker Taft
Harvard University Science Center
1 Oxford St
Cambridge MA 02138

Anyone exposed to small computer systems has used a language interpreter of some sort, and certainly may have thought about implementing their own interpreter. Unhappily, implementing an interpreter for a complete version of most computer languages is a difficult and time-consuming job, unsuitable for a part-time personal computer enthusiast. The language LISP provides a unique opportunity in this respect. The foundation for a very complete interpreter can be programmed by a single person in several months of part-time effort. As a bonus, the resulting interpreter provides the user with a high level language in which to express algorithms.

The Language

From the user's point of view, the primary data structure in LISP is the *list*. Every element of a list is either an *atom* or another list. An atom is a primitive named object, the name being an arbitrary string of characters:

ABC is an atom.
135 is an atom.
(ABC 135) is a list of two elements, both atoms.
((ABC 135) XYZ) is a list of two elements, the first of which is a list, the second is an atom.
(()) is a list of two elements, both being lists of zero elements. A list of zero elements, the *null* list, is identified with the atom NIL.

The feature of the language LISP which makes it at the same time a uniquely interesting language, and relatively easy to implement, is that all program elements are represented using these same kinds of objects: atoms and list. Constants, variables, expressions, conditionals, even function definitions are all represented using only atoms and lists.

A value is associated with each atom, allowing atoms to represent program variables and constants. A symbolic atom, like XYZ, would represent a variable. A numeric atom, like 237, would represent a constant.

Operations on variables and constants, like addition, or a function call, are represented by list expressions:

(ADD 2 5) would represent the expression $2 + 5$.
(SIN (MUL 2 Y)) would represent the expression $\sin(2y)$.

Conditionals, loops, and function definitions are also represented by list expressions, as illustrated by this recursive function implementing Euclid's greatest common divisor algorithm:

```
(DEF GCD (LAMBDA (X Y)
  (COND
    ((GREATER X Y) (GCD (SUB X Y) Y))
    ((GREATER Y X) (GCD X (SUB Y X)))
    (T X)
  )
))
```

This would be equivalent to the Pascal program:

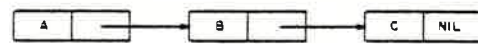
```
function gcd(x,y:integer):integer
begin
  if x>y then gcd := gcd(x-y, y)
  else
    if y>x then gcd := gcd(x, y-x)
  else
    gcd := x
end.
```

An important difference to note in the above comparison is that no explicit assignment to a function return value is made in LISP, whereas in Pascal one must explicitly say $gcd := \dots$ to specify the return value. In Pascal, and most other *procedural* languages, a distinction is made between program statements and expressions. In such languages some program statement must be executed to specify the return value, usually either a

The author(s) of the programs provided within have carefully reviewed them to ensure their performance in accordance with the specifications described. The author(s) and publisher however, make no warranties whatever concerning the programs and assume no responsibility or liability of any kind for errors in the program or for the consequences of any such errors. The programs are the sole property of the author(s).

Copyright 1979 BYTE Publications Inc. All Rights Reserved. BYTE and PAPERBYTE are registered Trademarks of BYTE Publications Inc. No part of this document may be translated or reproduced in any form without prior written consent from BYTE Publications Inc.

(A B C) IS BUILT UP OUT OF THREE DOTTED PAIRS



(J (K L M) N) IS BUILT UP OUT OF SIX DOTTED PAIRS

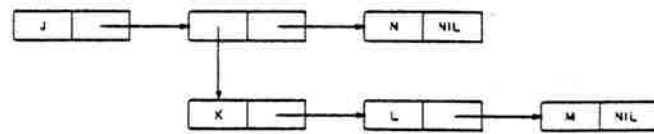


Figure 1: In most LISP systems, lists are built up out of dotted pairs which are two address cells. The left cell points to the first element of a list, and the right cell points to the rest of the list. The letters in the figure stand for atoms. NIL is a special atom used to signify the end of a chain of dotted pairs.

return statement or an assignment to the function name. In LISP, and other applicative languages, no such distinction is made. A function is simply a single expression, whose value is the return value of the subprogram.

This is made possible by built-in functions like COND used above. COND takes a list of two element lists as argument. It goes down the list of pairs, evaluating the first element of each pair. If the result is true (the atom T), the result of the entire COND is the value of the second element of the pair. If the value of the first element of the pair is false (the atom NIL), COND proceeds to the next pair. If COND reaches the end of the list, the result of the entire COND is simply NIL. In the above example this would never happen because the first element of the last pair is the atom T (whose value is always guaranteed to be itself, the atom T). This is the normal technique in LISP for using the COND function.

The expression:

```
(DEF GCD (LAMBDA (X Y)...
```

defines the atom GCD to be a function (or lambda expression) taking two arguments, to be called X and Y in the body of the definition. Notice that no explicit specification of the type of X or Y is provided. In LISP any arbitrary value, atom, or list may be the value associated with an atom. In this sense LISP is a typeless language. In fact the type of a value (ie: whether it is an atom or a list) is always determinable at execution time. Functions must check the types of the values of atoms if only certain types are legal arguments. In the above example the calls on GREATER and SUB would fail if the values associated with X and Y were not numeric atoms.

CARs and CDRs

Thus far we have only shown how to re-express algorithms written in a more conventional language, in the language LISP. The real power of LISP comes from its ability to directly manipulate lists, a data type not normally accessible in other languages. Three primitives, CAR, CDR (pronounced could-er), and CONS are provided for list manipulation. The function CAR takes a list

as argument, and returns the first element of the list, which may either be an atom or another list. The function CDR takes a list as argument, and returns the tail of the list, that is, all but the first element of the original list, as a new list. The function CONS takes two arguments, a new first element, and the tail of a list, and reconstructs a list, now one element longer. For example:

Assume the atom X is associated with the value:

(A B C)

Assume the atom Y is associated with the value:

(THE CAT IN THE HAT)

(CAR X) would be the atom A.

(CDR Y) would be the list (CAT IN THE HAT).

(CONS (CAR X) (CDR Y)) would be the list:

(A CAT IN THE HAT)

(CAR (CDR X)) would be the atom B.

In general the CAR of the CDR of a list is its second element, and a function called CADR is frequently defined as a kind of shorthand for CAR of the CDR.

You might wonder what would result if you gave two atoms as arguments to CONS, rather than an atom and a list. In most LISP systems this is in fact legal. The result reveals the underlying representation used for lists in LISP. In virtually all LISP systems, lists are built up out of dotted pairs, two-address cells, the left cell pointing to the first element of a list, and the right cell pointing to the rest of the list. This can be diagrammed schematically as in figure 1.

Because dotted pairs are used this way to build up lists, it is natural to call the left cell of a dotted pair the CAR and the right cell the CDR. (In fact the genealogy of the words CAR and CDR runs the other way. Dotted pairs were used in the initial implementation of LISP, and CAR and CDR referred to the address field and the decrement field of a word on the IBM 704.) Now you can perhaps guess that when you pass two atoms as arguments to CONS, you simply get a dotted pair with an atom in both the CAR and CDR. For example:



would be printed as:

(A . B)

The notation (A . B) is used whenever the CDR of the last dotted pair forming a linked list is a non-NIL atom. In general (D E F . NIL) would be equivalent to (D E F), whereas (D E F . G) could not be expressed without the dot notation.

Given the three primitives CAR, CDR, and CONS, and understanding the underlying representation of lists using dotted pairs, it is possible to write powerful list-manipulating programs in LISP. For example, suppose it is desirable to edit a large data structure, and change all occurrences of the symbol APPLE to ORANGE. In LISP we could easily write a routine called REPLACE which, given the data structure (ie: list structure), the original symbol (the atom APPLE), and the replacement symbol (the atom ORANGE), would go through the structure

and do the replacement, using itself recursively to do the replacement in all sublists of the list structure:

```
(DEF REPLACE (LAMBDA (STRUC OLD NEW)
  (COND
    ((EQ STRUC OLD) NEW)
    ((ATOM STRUC) STRUC)
    (T (CONS
        (REPLACE (CAR STRUC) OLD NEW)
        (REPLACE (CDR STRUC) OLD NEW)
      )
    )
  ))
```

Notice how the first two lines of the COND allow for the possibility that the input data structure is simply an atom (which may or may not be equal to the atom to be replaced). In addition, notice that the entire body of this function definition is a single COND, just as it was in the GCD example given above. This is frequently true in LISP programs. Finally, notice how the function simply passes the buck to recursive calls on itself if the STRUC argument is not an atom, CONSing together the results of the two inner calls. The reader is encouraged to go through an example of the execution of this function when the argument OLD is the atom APPLE, the argument NEW is the atom ORANGE, and the argument STRUC is the list structure:

```
(AN (APPLE A DAY) KEEPS (THE (APPLE MAN)
  BUSY))
```

The result should be:

```
(AN (ORANGE A DAY) KEEPS (THE (ORANGE
  MAN) BUSY))
```

If STRUC were:

```
(PEAR BANANA . APPLE)
```

the result should be:

```
(PEAR BANANA . ORANGE)
```

Other kinds of list-manipulating programs which are relatively easy to write in LISP, but very difficult in more conventional languages, include formula manipulation programs which might take in the list representation for a function (eg: (SIN (MUL 2 X))), and return the list representation for its derivative according to the rules of the calculus (eg: (MUL 2 (COS (MUL 2 X))))).

The author's system is being used for the development of a compiler/interpreter system which generates the list representation for a program written in a programming language, and then either interprets it directly, or generates the list of machine language statements to implement the program on a particular microcomputer. LISP makes such an undertaking quite straightforward (although not trivial, unfortunately!).

LISP Interpreter

Because programs are data objects (list structures) in LISP, the same routines used to read and print data objects may be used to read and print programs. Furthermore user functions, like a general list editor, can be used also to edit programs. This uniformity vastly simplifies the task of writing an interpreter for LISP. Only three basic modules need be produced: READ, EVAL, and

PRINT. READ accepts a LISP list expression from the terminal, in full parenthesized notation, and builds the internal representation of the list, sometimes called a form. EVAL takes a form as its single argument, and evaluates the form according to the LISP convention that the first element of such a list specifies the function, with the rest of the list as arguments.

The result of EVAL is another form. (The term form is sometimes reserved for LISP expressions which are legal input to EVAL. The term S-expression covers all types of lists, whether or not the first element is a legal function name. Within this paper, form will be used to refer to the internal representation of any type of LISP expression.)

PRINT takes a form as its argument, and types it on the terminal in fully parenthesized form. The top level loop of the LISP interpreter simply prompts the user for input (-> is the LISP prompt), READs in the user's input, EVALs the resulting form, and PRINTs the result of EVAL. In a conventional high level language syntax, this would be:

```
while true do begin
  patom("->");
  form := read( );
  form := eval(form);
  print(form)
end.
```

or in M6800 assembly language:

```
BIGLUP LDX PRMPAT  get prompt atom
        JSR  PATOM  print the atom
        JSR  READ   read the form typed in
* result now in M6800 x-register
        JSR  EVAL   eval the form
* result of EVAL back in x-register
        JSR  PRINT  print the form
        BRA  BIGLUP  and loop around
```

PATOM is a subroutine, also called by PRINT, when a form is known to be an atom. In an assembly language implementation, it would be very convenient to pass forms in the M6800 index (X) register. This register is 16 bits long, so it requires that forms be only 16 bits. Some representation must be chosen for all LISP objects so that a single 16 bit number may uniquely specify any arbitrary object. Dotted pairs are used to represent lists. Dotted pairs hold two forms, a CAR and a CDR, so they must be 32 bit objects. A natural choice is to allocate 4 consecutive M6800 memory bytes for dotted pairs, and specify dotted pairs by the address of their first byte. This means that any two different dotted pairs will be easily differentiated by the forms that specify them.

This still leaves the problem of deciding on an internal representation for atoms, including symbolic atoms, numeric atoms, and NIL. In the author's LISP system only two items of information are needed for each symbolic atom, the string of characters which are the print name of the atom, and the value currently associated with the atom (which is an arbitrary form). Again a 4 byte representation is chosen, with the first two bytes used as a memory address pointing to the first character of the print name, and the third and fourth bytes used to hold the value (a form) of the atom. Now the address of

this 4 byte object can specify the atom uniquely from all other atoms and from all other dotted pairs.

Unfortunately this does not provide a simple way of distinguishing atoms from dotted pairs, when just given the form. Several solutions to this problem are possible. One is to restrict dotted pairs to a certain part of memory, then the address would determine whether the form specified an atom or a dotted pair. A second method is to add an additional byte to both dotted pairs and atoms which simply contains a type specifier, say 1 for dotted pairs and 2 for atoms. This method makes future expansion of types simple, but is somewhat wasteful in terms of space. The third method, the one chosen for the author's system, is to align all dotted pairs and atoms on 4 byte boundaries, that is, with addresses which are a multiple of four. This means that the low order two bits of the address are expected to always be zero, and hence may be used to encode type information. In the author's system, dotted pairs are specified by forms with both bits zero, and symbolic atoms by 01 in the lower two bits. One of the bits is still unused, but will become very handy when *garbage collection* methods are discussed below.

With numeric atoms, their name determines their value, and hence only their name (or their value) need be specified by a form. On the author's M6800 system only hexadecimal memory addresses 0000 thru 7FFF were accessible for storage of dotted pairs and atoms, meaning that the high order bit of forms specifying either of these was always zero. A representation for numeric atoms was chosen to be a form with the high order bit set, 14 bits of numeric value, and one bit left for *garbage collection*.

A special representation for the NIL atom is used both because the value of NIL is, like numeric atoms, required always to be the atom itself, and because it is used universally to represent the end of a list. The form chosen to specify NIL is simply the value zero. In fact any form with the high order byte zero is treated like NIL to simplify the test for NIL in certain cases. This means that the 256 byte page starting at zero is not usable for storing atoms or dotted pairs, but this restriction causes no problem at all, since both are allocated starting at the highest address available, and the allocator runs into program long before it reaches page zero.

When writing a LISP interpreter, the implementor must decide relatively early on how forms will specify all types of LISP objects. Unfortunately, it may not be until well into the implementation that the implementor discovers that certain choices were inefficient or inconvenient.

One important requirement affecting this decision not yet mentioned is the need to implement the LISP EQ function. This function takes two arbitrary forms, and returns the atom T or the atom NIL depending on whether the forms specify the same dotted pair, or whether the forms specify the same atom. Whenever an atom is input by READ, it must return the form specifying that atom to the caller. Whenever the same symbolic atom name is typed, READ must return the same form, ie: a pointer to the same 4 byte cell. This is accomplished by retaining a linked list of all defined symbolic atoms (called the OBLIST).

Before allocating a new 4 byte cell for an atom, READ scans the OBLIST for an atom of the given print name. If found, READ returns a form specifying that pre-existing atom. (Otherwise it must copy the name into some area used for storing names, allocate a 4 byte cell, initialize the left cell to point to the name, and the right cell to NIL, and return a form specifying the new atom.) This method guarantees that two forms specify the same symbolic atom if and only if they have the same address.

In some implementations of numeric atoms, this same rule cannot be guaranteed. In such systems, numeric atoms are simply allocated an appropriately large cell to store their numeric value (and hence allowing numeric atoms greater than 14 bits), a new cell being allocated every time a new number is generated (which happens at every ADD, MUL, etc). In these systems it would be impractical to scan a list like the OBLIST every time any arithmetic calculation is done, and so the LISP function EQ may not rely on the rule that unequal forms indicate unequal atoms. In such systems, EQ must look at the contents of the cell specified by a numeric atom form, and make the comparison that way. In systems like the author's, EQ simply compares the forms themselves, no matter what type of atom the form may specify.

The choices made in representing the various types of LISP objects can be summarized in the high level language (Pascal-like) data structure specification in listing 1.

```

type lisptype =
  (dtpertype, symatmtype, numatmtype, nilatmtype);
dtptr =
  record
    car: form;
    cdr: form;
  end;
symatm =
  record
    name: 1array [0..n] of char;
    value: form;
  end;
form =
  packed record
    gcbit: boolean;
    case objtype: lisptype of
      dtpertype: (dtptrform: 1dtptr);
      symatmtype: (symatmform: 1symatm);
      numatmtype: (numatmform: -5000..4999);
      nilatmtype: ();
    end.

```

Listing 1: A Pascal data structure specification that could be used to represent various types of LISP objects.

READ Function

READ is the basic input routine for the LISP interpreter. READ accepts a fully parenthesized expression from the terminal, and builds up the internal representation, allocating new dotted pairs and atoms as necessary. If the expression is a list, READ returns a form specifying the first dotted pair of the constructed list. If the expression typed in is simply an atom, READ returns a form specifying the atom.

The logic of the READ routine is straightforward because the syntax of LISP expressions is so simple. READ calls a function RATOM to return the next input atom. RATOM actually does the work of allocating new 4 byte cells for symbolic atoms (when necessary) as explained above. RATOM returns a form specifying the

Listing 2: LISP (a) and M6800 assembly (b) code for the READ function.

```

(a) (DEF READ (LAMBDA (F)
      (READR (RATCM))
    ))
    (DEF READR (LAMBDA (A)
      (COND
        ((EQ A LPAREN) (READL (READ)))
        (T A)
      ))
    )
    (DEF READL (LAMBDA (F)
      (COND
        ((EQ F RPAREN) NIL)
        (T (CONS F (READL (READ))))
      ))
    )
    (b) READ JSR RATOM pick up the next input atom
          CPX LPARAT is it equal to the "(" atom?
          ONE RDRET no, simply return the atom
          JSR LSTINI yes, initialize a linked list
          RDLUP RDR READ call READ recursively to pick up next form
          CPX RPAREAT is it equal to the ")" atom?
          REC RDLUN yes, finalize the list and return
          JSR LSTADD no, allocate a new dotted pair,
          fill in the CAR, and add it to the list,
          and loop around.
          RDLUN JSR LSTEND finalize the list, get a pointer to
          the first dotted pair of the list
          and return.
          RDRET RTS

          * set up the stack for LSTADD.
          * first stack a NIL, then a pointer to the NIL.
          LSTINI LDX ZERO stack a NIL form
          JSR PUSHX
          LDX STKPTR stack a pointer to the NIL
          DEX as though it were the CDR of a cell.
          DEX
          JMP PUSHX push it and return.

          * add form in X-reg to list pointed to by value on top of stack
          * clobbers X-reg, A-reg, and B-reg.
          LSTADD JSR GETCEL get a dotted pair and fill in the CAR
          STX CELPTR save address of new cell
          JSR TOPX retrieve pointer from top of stack
          LDAA CELPTR link new cell onto list
          STAA CDR, X
          LDAD CELPTR+1
          STAB CDR+1, X
          LDX STKPTR update list pointer
          STAA CAR, X (which is on top of stack)
          STAB CAR+1, X
          RTS and return.

          * pop off list end pointer
          * return pointer to first dotted pair of list in X-reg
          LSTEND JSR POPX pop off and ignore list end pointer
          JMP POPX pop off and return list begin pointer.

```

Listing 4: LSTINI, LSTADD, and LSTEND build up a linked list of dotted pairs using two pointers on a stack, one to the first dotted pair, one to the dotted pair at the current end of the linked list.

Listing 3: M6800 implementation of linked-list stack manipulation routines.

```

* PUSHX saves the value of the X-reg on a linked-list stack.
PUSHX BSR GETCEL allocate a new dotted pair and fill in the CAR
      LDAA STKPTR fill in the CDR from STKPTR and fill in the CAR
      STAA CDR, X (CAR and CDR are symbols defined as
      LDAA STKPTR+1 0 and 2 respectively)
      STAA CDR+1, X
      STX STKPTR update the STKPTR
      LDX CAR, X restore the X-reg
      RTS and return.

* allocate a new dotted pair, save X-reg in the CAR,
* and set the CDR to NIL
GETCEL STX XTMP save the X-reg temporarily
      LDX FREPTR pick up a free dotted pair off the FREE list.
      DEC RCFREE no more left, so it goes...
      LDAA CDR, X update the FREE list pointer
      STAA FREPTR
      LDAA CDR+1, X
      STAA FREPTR+1
      CLR CDR, X set the CDR to NIL
      LDAA XTMP, X fill in the CAR from the X-reg
      STAA CAR, X
      LDAA XTMP+1
      STAA CAR+1, X
      RTS and return.

* restore the X-reg from the linked-list stack
POPX LIX STKPTR point to the top of the stack
      LDAA CDR, X update the STKPTR
      STAA STKPTR
      LDAA CDR+1, X
      STAA STKPTR+1
      BNA FRECEL and free up the dotted pair used.

* free up a dotted pair, and load X-reg from the CAR
FRECCEL STAA CDR, X
      LDAA FREPTR+1
      STAA CDR+1, X
      LDX CAR, X load X-reg from CAR
      RTS and return.

* return value at top of stack in X-reg
* (but leave it on stack)
TOPX LDX STKPTR return CAR of cell on top of stack.
      LDX CAR, X
      RTS

```

Listing 5: RATOM accepts characters one at a time from the terminal and builds them into atoms.

```

* return form in X-reg for next atom typed in
RATOM LDX SPCPTR save pointer to first open spot in name area
      LDX HAMPTR
      LDA #120
      BGT CHA
      BSR REALC
      BRA SPCPTR

* get next char of input, and store in global PEEKC
REALC JSR GETC use get char routine in ROM monitor
      STA PEEKC save it in PEEKC
      RTS and return with it also in A-reg.

* collect name of atom, build up form, etc.
GNAV BSR CCFYC save char in name area
      BSR SPCICH is it a special char ("(" or ")")?
      BSR RCLUP no, collect multi-character name
      BSR REALC yes, update PEEKC
      BSR SCAROL and go scan OBLIST
      BSR REATC get next character of name
      CPA #120 separator?
      BSR SCAROL yes, all done with name
      BSR SPCICH special char?
      BSR SCAROL yes, end atom now
      BSR CCFYC no, copy this char to name area
      BSR RCLUP and loop.

* copy char in A-reg to name area, and advance name space pointer
CGPYC LDX SPCPTR
      STA PEEKC
      INX store the char
      STX SPCPTR increment the pointer
      RTS and return.

* check if char is a special char (i.e. "(" or ")")
SPCICH CMA #("(" or ")")
      BEQ SPCLY 1. paren?
      BSR SPCLY or ". paren?
      RTS return with Z-bit set appropriately.

* null-terminate name, check if numeric, scan OBLIST if not.
SCANDI CLRA null-terminate the name
      BSR CCFYC check if numeric...
      BSR CRBYXR yes, go return numeric form
      BSR RETUM no, scan the OBLIST for the symbolic atom
      BSR OBLIST using LSTPTR to point to the elements of the list
      BSR LSTPTR end of list, must be new atom
      BSR HEVATM get next atom
      BSR CAR,X clear low-order bit of form
      BSR DEX get pointer to the atom's name
      BSR CAP,X compare with that at HAMPTR
      BSR CMJAM found it! return pre-existing atom
      BSR OLDATM not equal, go on to next atom
      BSR LSTPTR
      BSR CLR,X
      BSR SCANLP

* compare strings pointed to by X-reg and HAMPTR
* 7-bit set if equal
CMJAM STX XTMP? save X-reg temporarily
      LDX HAMPTR point at new atom name
      CMA #120 get next two chars
      LDA LDX
      LDA LDX

```

```

STX XTMP? save X-reg temporarily
LDX LDX point at other name
CMA #120 first char equal?
CMPRET nope, return with Z-bit clear
      BSR SPCPTR all done now?
      BSR CMJAM yes, return with Z-bit set
      BSR LDX second char equal?
      BSR CMPRET nope, return with Z-bit clear
      BSR STX all done now?
      BSR CMA yes, return with Z-bit set
      BSR RTS match so far, advance pointers

XTMP? and loop around.
XTMP Z-bit now set properly.
      BSR CMJAM
      BSR RTS

* we have an old atom, reset spcptr and return
OLDATM LDX LSTPTR get form specifying pre-existing atom
      LDX CAR,X
      BSR NILATM is this the atom with the print name "NIL"
      BSR RETUM yes, return a zero form instead
      BSR ZERC save form temporarily
      BSR FCRH reset SPCPTR
      BSR HAMPTR and return with FORM in X-reg
      BSR SPCPTR
      BSR FCRM
      BSR RTS

* we have a new atom, allocate a new cell and return new form
HEVATM LDX HAMPTR set up new atom cell with pointer to name
      JSR GETCEL in CAR of dotted pair
      INX set low order bit of form to indicate atom
      STX FORM save in FORM for later
      JSR GETCEL link new atom onto OBLIST
      LDA OBLIST
      STA CAR,X
      LDA OBLIST+1
      STA CAR+1,X
      BRA RATRET and return with FORM in X-reg

* here is a stripped down version of CKHUMB. It accepts
* only one-digit numeric atoms. Otherwise it
* returns with C-bit set.
CKHUMB LDX HAMPTR look at name
      O,X get first char
      SUBA #0 subtract out ASCII code for zero digit
      BLM not a digit if negative
      CMA #9 bigger than 9?
      BDM not a digit
      BDM more than one digit?
      BDM not allowed (for now)
      BDM #820 A-OK, build up numeric form in FORM
      BSR FCRH shift value around so that bit 1 is left open
      BSR DAB for garbage collector
      BSR FCRM+1 bits 1,2,3,4; bit 0 --> bit 0
      BSR FORM store low order byte of form
      BSR FORM and return it in X-reg
      BSR CLC with C-bit clear.
      BSR RTS bad number, return with C-bit set.
      BSR DABUM
      BSR SEC
      BSR RTS

```

```

(a) (DEF PRINT (LAMBDA (F)
      (PROGN
        (PRINT F)
        (PATOM NEWLINE)
        F
      )
    )
)
(DEF PRINR (LAMBDA (F)
  (COND
    ((ATOM F) (PATOM F))
    (T (PROGN
      (PATOM LPAREN)
      (PRINR (CAR F))
      (PRINL (CDR F))
      (PATOM RPAREN)
    )
  )
)
(DEF PRINL (LAMBDA (L)
  (COND
    ((DTPR L) (PROGN
      (PATOM SPACE)
      (PRINR (CAR L))
      (PRINL (CDR L))
    )
  )
)
)

(b) * type out a form, fully parenthesized, and then go to a new line.
PRINT JSR PUSHX save X-reg on stack
      BSR PRINR simply pass the buck to recursive PRINR
      LDX CRLFAT type out CR/LF
      BSR PATOM using PATOM
      JMP POPX restore X-reg and return.

* type out a form, with no CR/LF
* clobbers X-reg
PRINR JSR ISATOM is the form an atom?
      BCC PATOM yes, pass the buck to PATOM
      JSR PUSHX nope, stack the X-reg
      LDX LPARAT type out a "("
      BSR PATOM
      JSR TOPIX restore the X-reg
      LDX CAR,X type out the CAP
      BSR PRINR (recursively!)
      JSR POPX restore pointer to the dotted pair
      LDX CDR,X advance to next dotted pair in linked list
      JSR ISCTPR is there a next dotted pair?
      BCC PRPAR nope, go type a ")"
      JSR PUSHX yep, save the new X-reg again
      LDX SPACAT type out a space
      BSR PATOM
      BSR PRINL and loop around.
      BSR LDX type out a ")"
      BSR PRPAR and return (through PATOM).
      BSR RTS

```

Listing 6: LISP and M6800 assembly code of the PRINT routine.

atom typed. If this atom is anything but the atom "(" READ simply returns the atom as its result. If the atom returned by RATOM is "(", READ calls itself recursively until it gets the atom ")", meanwhile stringing the forms returned together as the CARs on a linked list of dotted pairs. This could be written as in listing 2.

In the LISP functions we are assuming that the atoms LPAREN and RPAREN were initialized to point to the atoms with print names "(" and ")" respectively. Notice that in the LISP version, READ accomplishes the loop of the machine code version with recursion in READL. The routines LSTINI, LSTADD, and LSTEND used in the assembly language version build up a linked list of dotted pairs, using two pointers on a stack, one to the first dotted pair, one to the dotted pair at the current end of the linked list. The pointers are on a stack so that READ may call itself recursively. The stack is actually a linked list itself. The linked-list stack is manipulated with the routines in listing 3. With these routines it is straightforward to implement LSTINI, LSTADD, and LSTEND for use in READ. These routines are shown in listing 4.

The primitive function RATOM turns out to be the real workhorse of READ. It is stuck with the job of accepting characters one at a time from the terminal, and building them up into an atom. RATOM must distinguish symbolic atoms from numeric atoms, and build up the corresponding forms. Atoms are in general separated by spaces, tabs, or carriage returns. However a few special characters always form single-character atoms when encountered (eg: "(" and ")") without any separator characters necessary.

In the author's LISP system RATOM is relatively sophisticated, allowing for atoms with spaces in their names if they are quoted ("..."). Also the single quote character ("'") is given special significance, as are "[]". However a simpler RATOM is quite enough for an initial implementation. To make this exposition simpler, only single digit numeric atoms will be allowed. Certainly in an eventual implementation, multidigit numeric atoms, optionally preceded by a minus sign would be accepted.

In this RATOM, the characters are copied into an area set aside to hold the names of atoms as they are input. A null character (ASCII code zero) is used to terminate the name, when a separator or special character is encountered. If the name is entirely numeric, then the atom is a numeric atom, and the form is simply the value of the number, with the high order bit set, and one other bit left zero for use in the garbage collector. Otherwise the atom is a symbolic atom, and a scan is made of the OBLIST for a pre-existing atom with the same name. If one is found, the characters just typed in are thrown away and a form specifying the pre-existing atom is returned. If the atom is a new one, a 4 byte cell is allocated (using GETCEL defined in listing 4) and a pointer to the new atom is added to the OBLIST. A form specifying the new atom is returned. The M6800 assembly language code for this is in listing 5.

PRINT Function

PRINT is the second major recursive function comprising the LISP interpreter. It takes a single form as argument, and types the value as a fully parenthesized LISP expression. PRINT simply calls the more primitive function PATOM when it is given an atom to type. Otherwise, PRINT types a left parenthesis, calls itself recursively to type out the elements of the list, and then types a right parenthesis. In any case, PRINT always types out a carriage-return/line-feed at the end. This can be coded as in listing 6.

In the LISP routines, the special function PROGN is used. PROGN simply evaluates all of its arguments in sequence, and then returns the value of the last one as the value of the entire PROGN. The two functions ATOM and DTPR are used to test the type of a LISP object. ATOM returns T if the argument evaluates to an atom — symbolic, numeric, or NIL. Otherwise ATOM returns NIL. DTPR is the exact opposite. It returns T if the argument evaluates to a dotted pair, and returns NIL otherwise. Such functions which return either T or NIL are called "predicates" in LISP in analogy with predicates as used in symbolic logic. Such functions in other languages are called Boolean functions.

Listing 7: A simplified version of PATOM which assumes single digit atoms.

```

* given atomic form in X-reg, type out name on terminal
* preserves X-reg
PATOM STX FORM save X-reg for later
      STX PHILAT go print 'NIL', if form is zero
      BEQ PHUMAT go print numeric atom (high bit set)
      RLT PHUMAT symbolic atom, clear low bit of form
      LDX CAR,X get pointer to print name of atom
      BRA PNAME and go type it out.

* NILM4 FCC 'NIL' string to print for NIL form
      FCU 0 (the null terminator)

* PHILAT LDA X point to a null-terminated string "NIL"
      DEC type out the chars one at a time
      JSR PATDUM until null char,
      INX using the ROM monitor put char routine.
      BNE PRAME advance to next character of name
          and loop around.

* PATDUM LDX FCR4 restore X-reg
      RTS and return.

* simplified version of numeric atom type-out
* form already stored in FORM. Only look at low order byte for now.
PHUMAT LDA FCR4+1 get low order byte, and shift it a bit
      LSR bits 2,3,4 --> 1,2,3
      AND #0 bit 0 --> bit 0
      ADD #10 add in ASCII code for zero
      CMPA more than a single digit?
      PHM2 yes, simply type out "" for now.
      PUTC send digit to monitor put char routine
      BRA PATDUM and return from PATOM.

```

Listing 8: A simplified version of EVAL.

```

* evaluate form passed in X-reg, return result in X-reg.
EVAL STX FORM save form temporarily
      BLE NUMNIL numeric atom or NIL, simply return
      LDA FORM+1 is this a symbolic atom?
      RORA (rotate bit 0 into carry bit)
      BCC no (low bit clear), must be dotted pair
      DEX yes (low bit set), load current value of atom
      LDX CDR,X
      STX FCR4 store it back in FCR4
      RTS and return with value in X-reg and FCR4.

* evaluate a function call (function specified by CAR of dotted pair)
EVLDTM LDX ALP save value of ALP, FLP, HLP on stack
      JSR PUSHY
      LEX FLP
      JSR PUSHX
      LDX HLP
      JSR PUSHX
      LDX FORM (now ALP, FLP, and HLP are "local" variables)
      BRA point back to original form

```

```

STX ALP save it in ALP (recursive call on EVAL
      LDX CAR,X will clobber FCR4)
      JSR EVAL evaluate CAR to get LAMBDA or SUBR expression
          (recursively!)
      BCC ISATOM is result an atom?
      JSR EVLERR yes, illegal expression!
      BCC EVLERR yes, illegal expression!
      (LAMBDA (...) ...
      or
      (SUBR . (carry-atom)
      STX FLP save pointer to function def
      LDX CAR,X now check the CAR for LAMBDA or SUBR
      JSR SUBRAT SUBR?
      BEQ EVLSUB yes, go evaluate args, etc.
      CPX LAMBDA?
      BEQ EVLLAM yes, go apply the LAMBDA
      JSR FERRMSG error break routine expects error message
      JSR ERRBRK error pointed to by X-reg.
      or error return, restore ALP, FLP, and HLP
      JSR FCR4 restore return value
      JSR POPX
      JSR POPX
      JSR POPX
      JSR ALP
      LDX FCR4
      RTS and return with result in X-reg
          and FCR4.

```

```

* evaluate the arg list, call the machine code subroutine.
EVLSUB JSR EVLALS EVAL list of args (pointer to it in FLP)
      LDX FLP point at CDR of function def
      JSR CDR,X
      JSR EVLERR error if numeric or NIL atom
      JSR TSATCM error if not atom at all
      JSR EVLERR point to print name of funny atom
      DEX
      LDX CAR,Y check for presence of special string:
      LDA 0,X (hex F21,400 used in this system)
      CMPA #F21
      BNE EVLERR
      LDA 1,X
      CMPA #400
      BNE EVLERR
      JSR FCR4 SUBR looks OK, call subroutine
      JSR FCR4 save returned result
      JSR FCR4 free up list of EVALed args
      LEX FORM restore X-reg
      BNE EVLXIT and return from EVAL.

* evaluate call on (LAMBDA (J K) ...)
EVLAM JSR EVLALS evaluate list of args
      JSR EVLRSV save old values of formals on stack
      LDX FLP store new values for formals
      LDX CAR,X point to body of LAMBDA
      JSR EVAL evaluate body (recursively!)
      JSR FLP save result form temporarily
      JSR FLP restore old values of formals
      LDX FLP restore X-reg
      BNE EVLXIT and return from EVAL.

```

EVAL Function

The EVAL function is the *heart* of the LISP interpreter. EVAL accepts one form as an argument, and evaluates it according to the LISP convention: the value of NIL is NIL, the value of a numeric atom is itself, the value of a symbolic atom is the form associated with the atom, and the value of a list is determined by applying the function specified by the CAR of the list to the list of arguments which make up the CDR of the list.

In most LISP systems at least two distinct kinds of functions exist, SUBRs and LAMBDA's. SUBRs are the built-in functions of the LISP system, written in machine code (like CAR, CDR, PATOM). LAMBDA's are the user-defined functions, defined like (DEF GCD (LAMBDA (X Y) ...)). The effect of such a DEF is simply to define the list (LAMBDA (X Y) ...) as the value associated with the atom GCD.

The type of object used to specify a SUBR function varies among LISP systems. Frequently a new type of ob-

ject is defined, called CODE, distinct from atoms and dotted pairs. A second alternative is to treat SUBRs like a funny kind of atom. The author's LISP system treats the bytes which make up the machine code of the SUBR like the print name of an atom. The SUBR is then specified by a dotted pair, with the CAR being the atom "SUBR" to identify the type of function, and the CDR being this atom with the funny print name. In fact the print name is prefixed with a special string which is unlikely to occur in a normal atom's print name, and hence PATOM could detect that the print name was not typeable, and simply type, say, "!" instead. In addition EVAL can check for the presence of this special string at the beginning of the print name to avoid treating a normal atom's print name as machine code. This method for specifying SUBRs avoids introducing an additional type, but the added complication in PATOM and EVAL may rule out the method in some implementations.

```

(DEF PRINL (LAMBDA (L)
  (COND
    ((DTPR L) (PROGN
      (PATOM SPACE)
      (PRINR (CAR L))
      (PRINL (CDR L))
    ))
    ((EQ L NIL) NIL)
    (T (PROGN
      (PATOM SPACE)
      (PATOM DOT)
      (PATOM SPACE)
      (PATOM L)
    ))
  ))
)

```

A corresponding change could be made to the assembly language routines.

As with the primitive function RATOM, the function PATOM turns out to be more difficult to implement than the recursive PRINT. PATOM must distinguish between symbolic atoms, numeric atoms, and NIL, and act accordingly. With symbolic atoms, PATOM simply types the null-terminated name of the atom. With numeric atoms, PATOM must convert back from the internal representation of the numeric value, to the string of ASCII characters which represent the number. With NIL, PATOM simply types 'NIL'. Listing 7 is a simplified version of PATOM with numeric atoms of only a single digit.

When EVAL is given a list to evaluate, it first evaluates the CAR of the list (recursively). The evaluation of the CAR should be either a LAMBDA expression, or a SUBR expression. If the evaluation of the CAR is an atom, or a list not headed by LAMBDA or SUBR, then EVAL stops, and indicates an error to the user.

If the CAR of the list gives a LAMBDA expression, the arguments to the function call are evaluated one at a time and saved on a list. The value associated with the "formal" arguments of the LAMBDA expression (eg: X and Y to the GCD routine given earlier) are saved on the stack. These formal arguments are then set one at a time to have the value of the corresponding actual arguments to the function (which were evaluated already). Finally, the "body" of the LAMBDA expression is evaluated, with the formal arguments now holding their new values. The result of evaluating the body is the result of the original function call. As a last step, EVAL restores the original values of the formal arguments.

Following the details of evaluation of such a function call is very difficult at first. The sequence of these steps is critical: evaluate actual arguments, save old values of formal arguments, set new values of formal arguments, evaluate body of LAMBDA, restore old values of formals. With any other sequence there is a chance that changes to the formal arguments of this function might interfere undesirably with the values of atoms in the calling routine's environment. These formal arguments are supposed to be strictly "local," that is, the choice of a name for a formal argument should be a strictly local decision, having no impact on variables with the same name in calling routines. Observing these rules allows LISP functions to be freely recursive. As the above examples of routines demonstrate, this recursion is in fact heavily used in LISP programming.

The steps in applying a SUBR function are simpler, because there are no formal arguments to worry about. EVAL simply evaluates the arguments to the SUBR, and passes them as a list to the machine code subroutine. EVAL expects the result of the SUBR to be left in register X when the subroutine returns.

This much of EVAL can be implemented on the M6800 as in listing 8.

The routines EVLALS, POPFRE, EVLNSV, EVLRSO, and EVLRST have not been included in listing 8 for brevity's sake. They are all relatively straightforward

routines, making heavy use of GETCEL, PUSHX, POPX, and FRECEL to build up and then release the lists of saved values.

Two additional types of LISP functions, normally recognized by an EVAL function, are called NLAMBDA and NSUBR (or FSUBR, or FEXPR if you prefer). These types of functions take their argument lists unevaluated. NSUBRs are simply passed the CDR of the original function call list, instead of a list of evaluated arguments. Similarly, NLAMBDA is provided with only a single argument, the list of unevaluated arguments. Without NSUBRs it is necessary for EVAL to recognize functions like COND as special cases, so that their argument list is not immediately evaluated. NSUBRs are specified in the same way as SUBRs, with the atom "NSUBR" replacing "SUBR" in the CAR of the dotted pair. PRINT will type out NSUBRs as "(NSUBR .!)"

NLAMBDA is very useful for creating elaborate user-defined functions which take argument lists that are as or more complicated than COND. NLAMBDA is necessary anytime the number of arguments is variable, or some of the arguments are wanted unevaluated.

To incorporate NLAMBDA and NSUBR in the above EVAL routines, two additional checks must be added immediately prior to EVLERR:

```

BEQ EVLLAM      ...
CPX NSUBAT      NSUBR?
BEQ EVLSU       yes, go call machine code
                subroutine
CPX NLAMAT      NLAMBDA?
BEQ EVLNLA      yes, pass list of args as single
                argument

```

* illegal exp...
EVLERR

and the additional routines EVLSU and EVLNLA must be included. Both of these routines are simpler than the corresponding routines EVLSUB and EVLLAM.

To make EVAL useful, some number of built-in SUBRs and NSUBRs must be written. The number of such built-in primitives can be kept quite small in LISP if they are chosen carefully. Most routines can be implemented as user functions if a few primitives exist. The primitives will certainly include PATOM, RATOM, EVAL, CAR, CDR, CONS, COND, SET, ADD, SUB, EQ, GREATER, ATOM, and NUMBER. All but SET and NUMBER have been used in the LISP function listings. SET is the primitive LISP assignment function. SET takes an atom and a value, and sets the value associated with the given atom to be the given value. NUMBER is a predicate function like ATOM, and simply returns T when its argument is a numeric atom. Listing 9 is an example of one of these primitives, the SUBR EQ.

Notice that the SUBRs and NSUBRs will start with the preface string (hex 21, 00 is used in this system). The argument list is always pointed to by ALP. Also notice

that the SUBR may not assume that the proper number of arguments were supplied. The general rule is to treat unspecified arguments as though they were NIL. In EQ above, this gives some rather strange behavior, where simply (EQ) will always return T. It still remains for the implementor to initialize the atom EQ to point to a dotted pair, (SUBR . funny-atom), with the print name of the funny atom set to point to the code at EQSBR as shown in listing 9. The final section of this article goes over some of the problems involved with this kind of initialization.

Garbage Collector

A garbage collector eventually becomes essential in any LISP system. It is possible to create dotted pairs that are no longer accessible to a LISP program by any path. This happens, for example, if a function like REPLACE is called and then the value returned simply PRINTed but not saved as a LISP atom. This cannot go on for long before all of the free space is used up with dotted pairs. The garbage collector's job is to find all of the dotted pairs.

The various algorithms for locating such jetsom of the LISP function evaluation process are all quite intricate. The basic idea is always to trace systematically down every list structure to its component atoms, marking every dotted pair encountered along the way. If a dotted pair is encountered which is already marked, then that branch of the list structure is assumed to be already fully traced. The garbage collector then makes a sequential scan of all of memory space occupied by dotted pairs, and links together all unmarked dotted pairs onto a special list, the free list. During the scan, the marked dotted pairs are simply skipped over, because they are assumed to still be a part of some useful list structure. When a marked dotted pair is skipped over, its mark is also cleared in anticipation of future garbage collections, when it might no longer be so lucky.

The difficulty with this trace and collect algorithm is that each dotted pair points to possibly two more dotted pairs, so during the tracing phase the garbage collector must eventually follow both paths. What this means

```

* two argument SUBR EQ
* return T if given identical forms, NIL otherwise
EQSBR  FCB  021  special preface string
        FCB  000
* ALP points to the list of evaluated arguments
LDX    ALP  get first arg
REQ    TRUE no args is equivalent to
        (EQ NIL NIL)
        which should return T.
        save first arg temporarily
        LDX  CAR,X
        STX XTMP
        LDX ALP  pick up second arg
        LDX CDR,X
        BEQ EQSNIL (EQ X) is equivalent to
        (EQ X NIL)
        LDX  CAR,X
        CPX  XTMP  are the forms identical?
        BEQ  TRUE  yes, return T.
        LDX  ZERO  no, return the NIL form
        RTS
*
TRUE   LDX  TATOM  return T atom
        RTS

```

Listing 9: EVAL may have built in primitives to expand the language. This is an example of the primitive SUBR EQ.

is that a second indication must be made on each dotted pair, indicating that the garbage collector is now busy tracing the CAR of this dotted pair, and will be returning later to trace the CDR of the dotted pair.

During the tracing phase, the garbage collector might very well be thought of as an ant determined to visit every branch of a tree. It goes out to the tip of each branch, but as it returns it must remember whether it has already traversed the other paths going out from each branching point. Even this analogy underrepresents the difficulty of a garbage collector, because the ant can simply turn around when it reaches the tip of a branch, but the garbage collector would normally have no clue as to how to climb back toward the root of a list structure once it gets out on a distant dotted pair.

The solution to the garbage collector's problem is to either reverse all the pointers in the list structure as it forays out to the terminating atoms and then reset the pointers on the way back in, or to keep a list of all dotted pairs which still require that their CDRs be traced. The first solution is like stringing a spool of thread behind you as you venture into an unexplored cave, following the thread back toward the mouth of the cave when you reach a dead end. Of course the same danger exists; that the delicate thread leading you back to the starting point might get tangled or broken.

The second solution is simpler, but suffers from the grave problem that it requires room to store the list of partially visited dotted pairs, and garbage collectors tend to be called upon at times when there is no more room to spare. In fact, the list of partially visited pairs need get no longer than the maximum "depth" of any list structure in the system, so that by setting aside a small portion of memory reserved for the use of the garbage collector's list, the implementor can get by with coding a much simpler tracing algorithm.

The author's system uses the pointer reversal method, and he will testify to the unlimited number of obscure problems which can appear during the debugging phase of its implementation.

It should be clear now why it was important to leave one bit in each form, and hence two bits per dotted pair, free for the use of the garbage collector. The bit in the CAR form can be used to indicate that the dotted pair has been visited once, and the bit in the CDR can be used to indicate that both paths from the dotted pair have been traced. These bits are only used during garbage collection, but because the garbage collector may be called at any time when GETCEL finds that there are no more 4 byte cells on the free list it may, in fact, run at almost any moment.

Because of this unpredictability, a LISP system with a garbage collector must be coded "defensively," jealously protecting any dotted pair allocated but not yet added to some accessible list structure. The machine code routines given in the listings do not all adhere to this rule. The reason for ignoring the garbage collector in the development thus far was simply to keep the design of the routines simple and relatively intuitive.

If the reader intends to include a garbage collector in an implementation of a LISP interpreter, more care must be taken. For example, two versions of the routine PUSHX would be defined. normal PUSHX and PROPSH

(protected push). The PROPSH would be used when the 16 bit value being pushed on the stack pointed to list structure which might not be accessible in any other way, and hence might get collected in the next garbage collection scan. PROPSH avoids this danger by marking the cell used to store the saved value so that the garbage collector will know to trace this form and its descendents.

Initialization

It is ironic, but somehow appropriate, that the section on initialization comes at the end of this article. Frequently it is in fact one of the last things an implementor thinks about. That is probably because initialization is one of the biggest difficulties facing the implementor of any language: assembler, interpreter, or compiler. By initialization is meant the inevitably awkward methods of getting the symbol tables, or the OBLIST in LISP preloaded with the names which are to be built-in to the system. Most of the routines written to enter symbols into symbol tables, or to add new atoms to the OBLIST, are all oriented toward names entered by the user of the language processor. The initialization phase of the system becomes quite complicated because of this orientation. The methods finally chosen are, in general, tedious, requiring a lot of special preparation by the writer of the initialization routine.

The best way to avoid these initialization difficulties is to spend a little extra effort in designing a few nice routines for taking information out of tables which are convenient for the implementor to set up and modify, and let these routines do the intricate bit-tiddling work necessary to get the objects in shape for the symbol table, or the OBLIST.

In the author's LISP initialization module are routines to build up dotted pairs in the form required for SUBRs and NSUBRs, and routines to allocate 4 byte cells for built-in atoms. The atom initialization routines are given the address of a contiguous table of null-terminated ASCII names, each followed by the address of a memory cell where the form specifying the new atom should be stored. This is where the symbols like TATOM, SUBRAT, LAMBAT, etc came from. They refer to memory locations in the base page of the M6800 (0 thru 255), where the forms specifying the atom T, SUBR, and LAMBDA, etc, are stored. The table to initialize these atoms was simply:

```

ATMTAB FCC  'T'
        FCB  0
        FDB  TATOM
        FCC  'SUBR'
        FCB  0
        FDB  SUBRAT
        FCC  'LAMBDA'
        FCB  0
        FDB  LAMBAT
        FCC  ...
.
.
.
        FDB  ...
        FCB  0  null-name terminates table

```

Although writing the special initialization routines was initially time-consuming, it was more than compensated for by the ease of adding more built-in atoms as the system grew.

Conclusion

We have traced through the implementation of a LISP interpreter and looked at a specific example for the M6800 processor. For further information on the garbage collecting routines and a complete listing of the interpreter, see listing 10.

Listing 10: The entire LISP interpreter for the M6800 with built in garbage collection.

```

0200          NAME LISP
0000          OPT NOG
0002          *SIMPLE LISP INTERPRETER
0004          OREGAD EQU $200
0006          CAR EQU 0
0008          CDR EQU 2
0010          EDI EQU 4
0012          * DIRECT PAGE STORAGE CELLS
0014          ORG $20
0016          BEGPR FDB BEGARR THIS SHOULD BE IN EACH MODULE
0018          STR1 RMB 2
0020          STR2 RMB 2
0022          STR3 RMB 2
0024          STR4 RMB 2
0026          STR5 RMB 2
0028          XTMP EQU STP1
0030          XTMP2 EQU STP2
0032          FREPR RMB 2
0034          ENDFR RMB 2
0036          SYHLST RMB 2
0038          NAMPTR RMB 2
0040          FORN RMB 2
0042          LSTPTR RMB 2
0044          SYMPTR RMB 2
0046          CELPR RMB 2
0048          SFSAVE RMB 2
0050          STMPTR RMB 2
0052          CNPTM2 RMB 2
0054          ALP RMB 2
0056          MLP RMB 2
0058          FLP RMB 2
0060          RLATH RMB 2
0062          LPARAT RMB 2
0064          RFARAT RMB 2
0066          DOTATH RMB 2
0068          SOUTAT RMB 2
0070          RUDATA RMB 2
0072          LANATH RMB 2
0074          MLANAT RMB 2
0076          SURRAT RMB 2
0078          ASURAT RMB 2
0080          TATON RMB 2
0082          CONATH RMB 2
0084          RSETAT RMB 2
0086          EDIATH RMB 2
0088          SURLS2 RMB 2
0090          LRKAT RMB 2
0092          RRKAT RMB 2
0094          CUREVL RMB 2
0096          SONPR RMB 2
0098          RDPFR RMB 2
0100          BCTEMP RMB 2
0102          GFTEMP RMB 2
0104          SIBLSSJ RMB 2
0106          * SINGLE CHAR SLOTS ***
0108          FEENC RMB 1
0110          RSIFLG RMB 1
0112          * GLOBAL CONSTANTS
0114          ORG $FC

```

```

00FC 6F A0      * END OF MEMORY, MINUS 192 FOR NUMFRM, FRMMON TABLES
00FE 00 00      * ZERO FOR $2000-192 MUST BE MULTIPLE OF 256, MINUS 192
0100          * GLOBAL VECTORS
0102          ORG $100
0104          E/VAL RMB 3
0106          FRJHT RMB 3
0108          GETCEL RMB 3
0110          FRECEL RMB 3
0112          POPX RMB 3
0114          TOPX RMB 3
0116          EVLATH RMB 3
0118          SETATH RMB 3
0120          LOOKUP RMB 3
0122          GNTXL RMB 3
0124          GFNTXL RMB 3
0126          LSTATH RMB 3
0128          LSTAD2 RMB 3
0130          LSTEND RMB 3
0132          LSTENO EQU POPX
0134          JNF ERKEX
0136          ATINI1 RMB 3
0138          ISDIFR RMB 3
0140          ISADOM RMB 3
0142          GETC RMB 3
0144          NUMJNV RMB 3
0146          NUMFRM RMB 3
0148          NUMNUM RMB 3
0150          PUTSYM RMB 3
0152          PUTC RMB 3
0154          JNP ERRBK
0156          BETSYM RMB 3
0158          STNFRG RMB 3
0160          GCOL RMB 3
0162          PROPSH RMB 3
0164          ISVAR RMB 3
0166          FRJNR RMB 3
0168          PROPOP RMB 3
0170          *
0200          ORG OREGAD
0202          *
0204          START EQU *
0206          STS SFSAVE
0208          LDX BEGPR
0210          STX SPCPTR
0212          LDX ENDREH
0214          STX ENDPTR
0216          *
0218          LRA A ENDREH SET UP HIGH BYTE OF XTMP
0220          STA A XTMP
0222          LDX M40 INITIALIZE TABLES TO 40'S
0224          LDA B M192
0226          EQU *
0228          STA A 0,X
0230          INX
0232          DEC B
0234          RNE NUMTLP
0236          * SET UP 64-BYTE TABLE
0238          LDX ENDREH

```

* INITIALIZE TABLES FOR NUMFRM AND FRMMON
* LAST 128 BYTES OF MEM HOLD TABLE TO CONVERT HIGH 7 BITS OF BCD TO 6 BITS OF BINARY
* PREVIOUS 64 BYTES HOLD TABLE TO CONVERT 6 BITS TO 7 BITS
* BIT # OF 128 BYTE TABLE IS ALWAYS ON

NUMBER OF FREE CELLS AFTER LAST GCOL INITIALIZED IN SURRSJ

0469 DE 32 LDX MANTPR
046B DF 2C SIX SPCPTR
* RETURN POINTER TO CELL
FOUND EQU *

0469 DE 34 LDX FORM
046F 9C 4A CFX NILATH
0471 26 04 BNE FND2
0473 DE FE LDX ZERO
0475 DF 34 FND1 EQU *

0477 FND2 EQU *

0477 39 RTS

0478
LKNH3 EQU *

* BUILD UP A NUMBER
* IS 16-BIT BCD NUMBER (+/-4999)
* HIGH BIT SET AS INDICATOR

0478 DE 32 LDX MANTPR
047A 4F CLR A
047B C0 20 SUB B B'-
047D 07 26 STA B XTMP
047F 27 11 REC LKNH5
0481 5F CLR B

0482
LKNH4 EQU *

* SHIFT NUMBER LEFT 4
* OVERFLOW?
* YEP-->ERROR MESSAGE

0482 B1 04 BHI NUMOVR
0484 22 23 ASL B
0486 58 ROL A
0487 49 ROL A
0488 58 ASL B
0489 49 ROL A
048A 58 ASL B
048B 49 ROL A
048C 58 ASL B
048D 49 ROL A
048E E0 00 ADD B B
0490 C0 30 SUB B B'-0

0492 08 INX
0493 60 00 TST
0495 26 EB BNE LKNH4

0497
LKNH6 EQU *

0497 97 28 STA A XTMP2
0499 07 29 STA B XTMP2+1
049B DE 28 LDX XTMP2
049D 96 24 LDA A XTMP
049F 26 03 BNE LKNH7

* YEP, DO A BCD TRICK

04A1 B0 01 45 JSR NUMINV
04A4 B0 01 48 JSR NUMFRN
04A7 20 80 BRA FNDIUM

* NUMOVR EQU *

04A9 DE 32 LDX MANTPR
04AB DF 2C SIX SPCPTR
04AD CE 05 59 LDX NUMOVR5
04B0 B0 01 54 JSR ERBRK
04B3 20 E0 BRA FND1

IXLST EQU *

* PTR TO LIST OF CHARS IN X-REG
* CHAR IN A-REG, RETURNS INDEX IN B-REG
* B < 0 IF NOT FOUND (CC S SET ON EXIT)
* SIX XTMP
* BSR INLIST
* BCC IXL2
* LDA B N-1
* NOPE

04B5
04B5 DF 26
04B7 B0 10 BSR
04B9 24 04 BCC
04BB CA FF LDA B N-1

04B8 20 06 BRA IXLRT
04BE DF 28 SIX XTMP2
04C1 B6 29 LDA B XTMP2+1
04C3 B0 27 SUB B XTMP+1
04C5 EQU *

04C5 DE 26 LDX XTMP
04C7 50 TST B
04C8 39 RTS

INLIST EQU *

* PTR TO CHAR LIST IN X-REG, CHAR IN A-REG
* RETURNS WITH CARRY SET IF NOT FOUND, CLEAR OTHERWISE
* X-REG CLOBBERED (LEFT POINTING AT MATCH)
* MATCH CHAR IN LIST

04C9 A1 00 CMA 0,X
04CB 27 06 BEQ INLRT
04CC 08 INX
04CE AD 00 TST 0,X
04D0 26 F7 BNE INLIST
04D2 00 SEC
04D3 EQU *

* YUP, SET CARRY
* RETURN WITH CARRY SET OR CLEAR

04D4
GETSYM EQU *

* READ NEXT SYMBOL, ADD TO SYNLST, RETURN PTR IN X-REG
* SPCPTR, ERBRK, SYNLST, ENDPTR MUST BE INITIALIZED
* USES GETC SUBROUTINE TO GET CHARS
* PEKCK SHOULD BE ZERO INITIALLY
* TYPE OF LEXEME RETURNED IN B-REG...
* 1 = QUOTED STRING
* 2 = BREAK CHAR (LINE ' ' OR '{')
* 3 = ALPHANUMERIC
** 4 = NON-ALPHANUMERIC ** REMOVED THIS CATEGORY 11/24/78 **
* <SP> AND ALL CTL CHARS (E.G. <HT> AND <FF>)
* ACT AS SEPARATORS, EXCEPT <EOI> IS BREAK

04D4 5F CLR B
04D5 96 F1 LDA A RSTFLG
04D7 26 38 BNE RETE12
04D9 DE 2C LDX SPCPTR
04DB DF 38 SIX SYMPTX
04DD 96 F0 LDA A PEKCK
04DF 26 03 BNE G52

04E1 B0 01 42 JSR GETC
04E4 EQU *

04E4 CE 05 51 LDX BRKRLST
04E7 80 E0 BSR INLIST
04E9 24 09 BCC G5BRK
04EB 81 20 CMA A
04ED 22 27 BHI G5NSEP
04EF 50 TST B
04F0 26 12 BNE ENDSYM
04F2 20 E0 BRA G5NXTC
04F4 EQU *

04F4 50 TST B
04F5 26 0D BNE ENDSYM
04F7 81 04 CMA A
04F9 27 13 BEQ RETE01
04FB 80 42 BSR COPYC
04FD 81 22 CMA A
04FF 27 24 BEQ G5OSTR
0501 C6 02 LDA B W2
0503 4F ENDPK EQU *

0504 4F ENDSYM EQU *

0504 97 F0 STA A PEKCK
0506 4F CLR A
0507 80 36 BSR COPYC
0509 DE 38 LDX SYMPTX

0508 7E 03 FC JMP LOOKUP AND RETURN DIRECTLY FROM LOOKUP

050E RETE01 EQU *

050E 7F 00 F0 CLR PEKCK

0511 RETE12 EQU *

0511 DE 64 LDX EGIATH
0513 DF 34 SIX FORM
0515 39 RTS

* G5NSEP EQU *

** NOT A SEPARATOR OR BREAK, CHECK IF OTHER SPECIAL
* LDX BSFCLST
* BSR INLIST
* BCC G5SPCL
* NOT SPECIAL, ASSUME "ALPHANUMERIC"
* FIRST CHAR?

0516
0516 50 TST B
0517 26 06 BNE G5ANCK
0519 CA 03 LDA B W3
051B EQU *

051B 88 22 BSR COPYC
051D 20 C2 BRA G5NXTC
051F EQU *

051F C1 03 CMA B W3
0521 27 FB BEQ G5ANST
0523 20 DF BRA ENDSYM

* G5SPCL EQU *

* TSTB FIRST SPECIAL CHAR?
* RNE G5SPCK
* LDR W4 YES, SET SPECIAL CHAR SYM CODE
* BRA G5ANST AND GO STORE CHAR
* G5SPCK EQU *

* CMA B W4 NOW STORING SPECIAL CHARS?
* BEQ G5ANST
* BRA ENDSYM
* NOPE, GO FINISH UP PREVIOUS SYMBOL
* QUOTED STRING TYPE SYMBOL

0525
0525 C6 01 LDX B W1
0527 EQU *

0527 80 01 42 JSR GETC
052A B0 13 BSR COPYC
052C 81 04 CMA A
052E 27 04 BEQ ENDSYM
0530 B1 22 CMA A
0532 26 F3 BNE G502
0534 B0 01 42 JSR GETC
0537 81 22 CMA A
0539 26 C9 BNE ENDSYM
053B B0 02 BSR COPYC
053D 20 E8 BRA G502

COPYC EQU *

* SAVE CHAR IN A-REG AT LOC PT'D BY SPCPTR
* CHECK AGAINST ENDPTR FOR SAFETY

053F DE 2C LDX SPCPTR
0541 9C 2E CPX ENDPTR
0543 2A 06 BPL CFCBAD
0545 A7 00 STA A 0,X
0547 08 INX
0548 DF 2C SIX SPCPTR
054A 39 RTS

CFCBAD EQU *

* NO MORE FREE SPACE
* JMP ERREX

054B CE 05 6A LDX NUMOVR5
054E 7E 01 36 JMP ERREX

7118 PSTRG EQU #7118
711E PCRLF EQU #711E
* TABLES, ETC.
BRKLT FCC
* (USED TO INCLUDE " " AND "()" IN ABOVE LIST)

0551 22
0556 27

0557 04 FCB EOI
0558 00 FCB 0

* SFCLST FCC 'IMBZ=***:~<?/!~>
* FCB 0

0559 4E NUMOVS FCC 'NUMERIC OVERFLOW'
0569 04 FCB 4
056A 4E NUMOVS FCC 'NO MORE FREE SPACE'
057C 04 FCB 4
057D EQU *

REGADR EQU *

END EQU *

NO ERRORS DETECTED

SYMBOL TABLE:

ALP 0044 ATHINI 0139 REGADR 057D REGPTR 0020 BRKLT 0551
CAR 0000 CDR 0002 CELPTR 003A CNFTH2 0042 CNFINP 0040
CONATH 0060 COPYC 053F CFCBAD 0548 CUREVL 006C DADPTR 0070

DOTATH 0050 ENDMEN 00FC ENDPK 0503 ENDPTR 002E ENDSYM 0504
EOI 0004 EOLATH 0064 EOL 0434 ERBRK 0154 ERREX 0136
EVAL 0100 EVALTH 0118 FLP 0048 FND1 0475 FND2 0477
FNDIUM 0459 FORM 0034 FOUND 046D FRECEL 010F FREPTR 002A
FNDIUM 0148 GCFREE 0074 GOOL 015D GCTEMP 0072 GETC 0142
GETCAR 0163 GETCEL 010C GETSYM 0404 GEXLTL 0127 GEXTL 0124
G52 04E4 G5ANCK 051F G5ANST 051B G5BRK 04F4 G5NSEP 0516
G5NXTC 04E1 G502 0527 G5OSTR 0525 INLIST 04C9 INLRT 0403
ISATON 013F ISDTR 013C ISE0 03F4 ISVAR 0166 IXL2 0418
IXLST 04B5 IXLRT 04C5 LANATH 0056 LBRAT 0068 LK01 0418
LK02 041C LK03 042E LK04 0458 LK05 0459 LKUNH2 0410
LKUNH3 0478 LKUNH4 0482 LKUNH5 0492 LKUNH6 0497 LKUNH7 04A4
LKUNH8 0408 LOOKUP 03FC LPARAT 004C LSTAD2 0130 LSTAD3 012B
NLATH 0115 LSTEAD 0133 LSTEHI 012A LSTPTR 0036 MANTPR 0032
NSUBAT 005C NUMFRM 0148 NUMINV 0145 NUMOVR 0459 NUMOVS 054A
PCRLF 711E PEKCK 00F0 POPX 0115 PRTHR 0169 PRINT 0106
PROP5H 0160 PSTRG 7118 PUSHX 0112 PUTC 0151 PUTSYM 014E
RPARAT 0054 RBRKAT 006A READ 0103 RETE12 0511 REIEDI 050E
RSETAT 004E RSTFLG 00F1 SETATH 011E SOMPTR 004E
SPCPTX 002C SFSAVE 003C SQUAT 0052 STE2 030A STKPTR 015A
STKPTR 003E STP1 0026 STP2 0028 STRI 0022 STR2 0024
STRED 0300 SUBLS2 0066 SUBLS3 0076 SUBRAT 005A SYNLST 0030
SYMPTX 0038 TATON 005E TOPX 0118 TYP5U1 0109 UNED 03FB
XTMP 0026 XTMP2 0028 ZERO 00FE

05A0 05A0 0582 ENDS LOOKUP
 0000 EQU 0
 0002 EQU 2
 0004 EQU 4
 * DIRECT PAGE STORAGE CELLS
 0020 EQU 120
 0020 06 8C BEGPR FDB BEGADR THIS SHOULD BE IN EACH MODULE
 0022 STR1 RAB 2
 0024 STR2 RAB 2
 0026 STP1 RAB 2
 0028 STP2 RAB 2
 002A XTMP EQU STP1
 002B XTMP2 EQU STP2
 002A FREPTR RAB 2
 002C SFCPTR RAB 2
 002E ENDPTR RAB 2
 0030 SYNLSI RAB 2
 0032 HAMPTR RAB 2
 0034 FORN RAB 2
 0036 CAPTMP RAB 2
 0042 ALP RAB 2
 0044 NLP RAB 2
 0046 FLP RAB 2
 0048 NILATM RAB 2
 004C LPARAJ RAB 2
 0050 POTATM RAB 2
 0052 SOUTAT RAB 2
 0054 BUOATH RAB 2
 0056 LANATH RAB 2
 0058 NLANAT RAB 2
 005A SUBRAT RAB 2
 005C ASUBAT RAB 2
 005E TATOM RAB 2
 0060 COMATH RAB 2
 0062 RGETAT RAB 2
 0064 EGTATM RAB 2
 0066 SUBLS2 RAB 2
 0068 LBRKAT RAB 2
 006A RBRKAT RAB 2
 006C COREVL RAB 2
 006E SOMPTR RAB 2
 0070 BDPTR RAB 2
 0072 GCTEMP RAB 2
 0074 BCFREE RAB 2
 0076 SUBLSI RAB 2
 * SINGLE CHAR SLOTS ...
 00F0 FEERC RAB 1
 00F1 ASIFLG RAB 1
 * GLOBAL CONSTANTS
 00FC EQU 0
 00FE EQU 0
 0100 EQU \$100
 0102 EQU 2
 0104 EQU 4
 0106 EQU 6
 0108 EQU 8
 010A EQU 10
 010C EQU 12
 * USED TO BE TYFSMT **

010F FRECEL RAB 3
 0112 PUSIX RAB 3
 0115 POPX RAB 3
 0118 TOPX RAB 3
 011B EVLATM RAB 3
 011E SETIATM RAB 3
 0121 LOOKUP RAB 3
 0124 GAXTL RAB 3
 0127 GFAXTL RAB 3
 012A LSTINI RAB 3
 012D LSTADD RAB 3
 0130 LSTAD2 RAB 3
 0133 LSTEND RAB 3
 0135 LSTEND EQU POPX
 0136 ERKEX RAB 3
 0139 ATHINI RAB 3
 013C ISDIPR RAB 3
 013F ISATOM RAB 3
 0142 BETA RAB 3
 0145 NUNINV RAB 3
 0148 7E 06 0B MUNFRM
 0148 7E 05 E9 JMP
 014E 7E 05 A0 JMP
 0151 PUTC RAB 3
 0154 ERBRK RAB 3
 0157 SETSYM RAB 3
 015A STMFRG RAB 3
 015D GCOL RAB 3
 0160 PROPSH RAB 3
 0163 GETCAR RAB 3
 0166 ISVAR RAB 3
 0169 7E 06 3F JMP
 016C PRADPOP RAB 3
 711B PSTRWG EQU \$711B
 711E FCRLF EQU \$711E
 *
 05A0 ORG OBEGAD
 05A0 PUTSYM EQU *
 * PUTSYM EXPECTS ATOM PTR IN X-REG
 * USES PUTC TO OUTPUT CHARS
 05A0 DF 34 STX FORN
 05A2 28 17 RNI PUTNUM NUMBER--> PUTNUM
 05A4 26 02 RNE PSY2
 05A6 DE 4A LDX NILATM
 05A8 EQU *
 05A8 09 REX
 05A9 EE 00 LDX CAR,X
 05AB 8D 03 BSR PUTSTR
 05AD DE 34 LDX FORN
 05AF 39 RTS
 05B0 PUTSTR EQU *
 * PUTSTR EXPECTS ADDR OF ZERO-TERMINATED CHAR STRING
 * IN X-REG (WHICH IS CLOBBERED!)
 * USES PUTC TO OUTPUT INDIVIDUAL CHARS.
 05B0 EQU *
 05B0 A6 00 LDX A 0,X
 05B2 27 0A RNE PSRT
 05B4 8D 01 51 JSR PUTC
 05B7 08 INX
 05B8 20 F6 BRA PSNXT
 05BA 39 EQU *
 05B8 PSRT EQU *
 *
 05B8 PUTNUM EQU *
 * PUTNUM EXPECTS FORM IN X-REG
 * USES PUTC TO OUTPUT IT

05B8 8D 2C BSR FORNUM
 05B8 2A 08 BFL PMH2
 05B8 86 2D LDA M *-
 05C1 8D 01 51 JSR PUTC
 05C4 8D 01 45 JSR NUNINV
 05C7 EQU *
 05C7 BF 2A STX XTMP
 05C9 8A 2A LDB XTMP
 05CB 96 27 LDA A XTMP+1
 05CC EQU *
 05CC 5D FSH B
 05CF 5D RNE PNMH2
 05D0 26 04 CMP A #409
 05D2 81 09 BLS PNMH3
 05D4 23 0E EQU *
 05D6 EQU *
 05D6 36 FSH A
 05D7 54 LSR B
 05D8 44 ROR A
 05D9 54 LSR B
 05DA 46 ROR A
 05DB 54 LSR B
 05DC 46 ROR A
 05DD 54 LSR B
 05DE 44 ROR A
 05DF 8D EE BSR PNMH2
 05E1 32 PUL A
 05E2 84 0F AND A #40F
 05E4 EQU *
 05E4 8A 30 ORA M *-
 05E6 7E 01 51 JMP PUTC
 05E9 FORNUM EQU *
 * EXPECTS FORM IN X-REG
 * RETURNS 4-DIGIT BCD NUMBER IN X-REG
 * USING 10'S COMPLEMENT
 * M-BIT SET IF NUMBER IS NEGATIVE
 * Z-BIT SET IF NUM IS ZERO
 05E9 DF 34 STX FORN
 05E9 96 FC LDA A ENDMEN
 05E9 97 24 STA A XTMP
 05EF 9A 35 LDA A FORN+1
 05F1 16 TAB
 05F2 00 SEC
 05F3 46 LSR A
 05F4 44 LSR A
 05F5 8D 30 ROR A
 05F7 54 LSR B
 05F8 49 ROR A
 05F9 97 29 STA A XTMP2+1
 05F9 96 34 LDA A FORN
 05FD 44 LSR A
 05FE 8D 27 BSR FNX
 0600 49 ROR A
 0601 97 20 STA A XTMP2
 0603 DE 28 LDX FORN
 0605 27 03 BEQ FNMKRT
 0607 88 30 ADD A #830
 0609 4D TST A
 060A EQU *
 060A 39 FNMKRT EQU *
 *
 0608 NUMFRM EQU *
 * EXPECTS BCD NUMBER IN X-REG (10'S COMPLEMENT)
 * CONVERTS TO FORM WITH HIGH BIT SET, AND BIT 1=0 FOR GCOL
 0608 BF 28 STX XTMP2
 0608 96 FC LDA A ENDMEN
 0609 97 26 STA A XTMP
 0611 96 29 LDA A XTMP2+1
 CONVERT LOW BYTE

0613 1A TAB
 0614 8D 0F BSR MFX
 0616 40 ASL A
 0617 36 ROR B
 0618 49 ROL A
 0619 97 35 STA A FORN+1
 061B 96 28 LDA A XTMP2
 061D 8D 06 BSR MFX
 061F 49 ROL A
 0620 97 34 STA A FORN
 0622 DE 34 LDX FORN
 0624 39 RTS
 *
 0625 EQU *
 0625 00 SEC
 0626 46 ROR A
 0627 EQU *
 0627 97 27 STA A XTMP+1
 0629 DE 24 LDX XTMP
 062B 8A 00 LDA A 0,X
 062D 39 RTS
 *
 062E 96 F1 LDA A RSTFLG
 0630 27 01 RNE PRINT1
 0632 EQU *
 0633 RTS
 0633 8D 0E BSR PRIMR
 0635 DF 34 STA A FORN
 0637 8D 00 LDA A #D
 0639 8D 4B BSR PPUTC
 063B 8A 0A LDA A #A
 063D 20 47 MRA PPUTC
 *
 063F EQU *
 063F 96 F1 LDA A RSTFLG
 0641 26 EF RNE PRINT
 0643 EQU *
 * PRINT TAKES ADDR OF OBJECT IN X-REG
 * ASSUMES EITHER DIPR, NIL, ATOM, OR NUMBER
 * CALLS PUTSYM AND PUTC
 * ASSUMES "NIL" ATOM STORED AT "NILATM"
 * USES PUSHX SUBR TO STACK X-REG
 * ALSO POPX AND TOPX TO REFERENCE STACK SWITCH ON TYPE OF FORN
 0643 8D 01 3F JSR ISATOM
 0646 25 03 BCS DIFAT
 0648 7E 05 A0 JMP PUTSYM
 *
 0648 EQU *
 0648 8D 01 60 JSR PROPSH
 064E 8D 01 5A JSR SIKFRG
 * GOT A DIPR, PRINT "(!"
 0651 86 28 LDA A #7!
 0653 8D 31 BSR PPUTC
 0655 EE 00 LDX CAR X
 0657 8D 0A BSR PRIMR
 0659 8D 01 18 JSR TOPX
 065C EQU *
 065C EE 02 LDX CAR X
 065E 8D 01 3C JSR ISDIPR
 0661 27 1A RNE PRT1
 0663 25 0E BCS PRT3
 0665 EQU *
 0665 8D 1D BSR PPUTSP
 PRINT A SPACE

SET UP LOW BIT
 STORE THE RESULT
 CONVERT THE HIGH BYTE
 SET UP THE LOW BIT
 RETURN THE RESULT
 LOW BIT--> C-BIT, 1--> HIGH BIT
 STORE LOW BYTE OF FORNUM/NUMFRM TABLE ADDR
 LOAD CONVERTED VALUE
 AND RETURN
 "PRINT" TYPES OUT FORM, FOLLOWED BY CR,LF
 USES RECURSIVE ROUTINE "PRINR" TO DO THE WORK
 FORM RETURNED IN X-REG AND "FORN"
 PRINT EQU *
 LDA A RSTFLG IN A RESET
 RNE PRINT1
 PRINT EQU *
 YES, RETURN IMMEDIATELY
 CALL RECURSIVE PRINT ROUTINE
 SAVE IN FORN
 CR,LF
 TYPE OUT FORM, WITHOUT CR,LF
 PRIMRO EQU *
 LDA A RSTFLG IN A RESET?
 RNE PRINT
 YES, SIMPLY RETURN
 PRINT EQU *
 PRINT TAKES ADDR OF OBJECT IN X-REG
 ASSUMES EITHER DIPR, NIL, ATOM, OR NUMBER
 CALLS PUTSYM AND PUTC
 ASSUMES "NIL" ATOM STORED AT "NILATM"
 USES PUSHX SUBR TO STACK X-REG
 ALSO POPX AND TOPX TO REFERENCE STACK SWITCH ON TYPE OF FORN
 JSR ISATOM
 BCS DIFAT
 JMP PUTSYM
 NIL, ATOM, OR NUMBER
 DIPRT EQU *
 JSR PROPSH
 JSR SIKFRG
 SAVE X-REG FOR RESTORE LATER
 FRAGMENT THE STACK
 GOT A DIPR, PRINT "(!"
 LDA A #7!
 BSR PPUTC
 LDX CAR X
 BSR PRIMR
 RESTORE X-REG!
 JSR TOPX
 AND FOLLOW LIST
 LDX CAR X
 JSR ISDIPR
 CHECK TYPE OF FORM
 RNE PRT1
 A NUMBER OR AN ADDR(!)
 BCS PRT3
 EQU *
 BSR PPUTSP
 PRINT A SPACE

0667 BD 01 I2 JSR PUSHX SAVE X-REG PRINT THE CAR
066A EE 00 LDA CAR,X PRINT THE CAR RECURSIVELY!
066C 8D 05 RSR FRJNR RESTORE X-REG
066E BD 01 I5 JSR POFX
0670 2E 01 BRN PRTZ
0673 PRTJ35 EDU *
* ATON OR NUMBER AT END OF LIST
* USE ".," SYNTAX PUT OUT ". "
* DIRECT PAGE STORAGE CELLS THIS SHOULD BE IN EACH MODULE

0673 8D 0F RSR PPUTSP
0675 86 2E LDA A M,"
0677 8D 00 RSR PPUTSP
0679 8D 09 RSR PPUTSP
067B 8D C6 RSR FRJNR
067D EDU *
0670 86 29 LDA A M,"
067F 8D 05 RSR PPUTSP
0681 7E 01 C6 JNF PPROPUP
* END THE LIST
* RESTORE X-REG AND RETURN
0684 PPUTSP EDU *
* PUT OUT A SPACE
068A 86 20 LDA A M,"
0686 EDU *
0686 7E 01 51 JNF PUTC
* USED TO AVOID LONG JSR
0689 EDU *
0689 3F SUI UNDEF
068A 20 FB BRA UNDEF
068C REGADR EDU *
* STOP, CALLED TYPSTW FROM SOMEWHERE

NO ERROR(S) DETECTED SYMBOL TABLE:

Table with 2 columns: Address and Symbol. Symbols include ALP, CDR, CUREVL, ENDPTR, EVAL, FORH, GCOL, GETSYM, ISVAR, LSTARB, LSTPR, NLP, PCRLF, PWHRR, PRIHR, PROFSH, PSXHT, PUTC, RBRKAT, SETATM, STKFRG, SYMPTR, XTMP2.

NO ERROR(S) DETECTED SYMBOL TABLE:

Table with 2 columns: Address and Symbol. Symbols include EVAL, HEAD, PRINT, JNP, FRECEL, PUSHX, POPX, TOPX, EVALM, SETATM, LOOKUP, GNTXL, GFXTL, LSTINI, LSTABD, LSTEND, LSTENO, ERXK, ATHNI, ISDIPR, ISAIOM, GETC, NUMIN, NUMFM, FRNMD, PUTC, ERKBRK, BEISYN, SINKFG, GCOL, PROFSH, GETCAR, ISVAR, PRINTX, FRADOP, ORG, UREGAD, SJR STKFRG, FRAGMENTING THE STACK.

* (XREG#1) CONTAINS OLD SP OF FORMER FRAGMENT
* (XREG#2) CONTAINS ADDR OF "UNRAVL" SUBROUTINE
* (XREG#0) CONTAINS RETURN ADDRESS FOR THIS SUBROUTINE
* CHECK LDW BYTE OF SP FOR MIN. STACKLIM
* SAVE X-REG FOR RESTORE LATER
* ADJUST ENDPTR TO 256-BYTE BOUNDARY
* FREE UP BLOCK DOWN TO 256-BYTE BOUNDARY
* UPDATE ENDPTR
* ALLOCATE NEW 256-BYTE STACK FRAGMENT
* BUMPING INTO NAME SPACE!
* NOPE, UPDATE ENDPTR
* LDA A W250

0684 86 20 EDU *
0686 7E 01 51 EDU *
0689 3F SUI UNDEF
068A 20 FB BRA UNDEF
068C REGADR EDU *
068E 86 20 EDU *
0690 86 20 EDU *
0692 86 20 EDU *
0694 86 20 EDU *
0696 86 20 EDU *
0698 86 20 EDU *
069A 86 20 EDU *
069C 86 20 EDU *
069E 86 20 EDU *
06A0 86 20 EDU *
06A2 86 20 EDU *
06A4 86 20 EDU *
06A6 86 20 EDU *
06A8 86 20 EDU *
06AA 86 20 EDU *
06AC 86 20 EDU *
06AE 86 20 EDU *
06B0 86 20 EDU *
06B2 86 20 EDU *
06B4 86 20 EDU *
06B6 86 20 EDU *
06B8 86 20 EDU *
06BA 86 20 EDU *
06BC 86 20 EDU *
06BE 86 20 EDU *
06C0 86 20 EDU *
06C2 86 20 EDU *
06C4 86 20 EDU *
06C6 86 20 EDU *
06C8 86 20 EDU *
06CA 86 20 EDU *
06CC 86 20 EDU *
06CE 86 20 EDU *
06D0 86 20 EDU *
06D2 86 20 EDU *
06D4 86 20 EDU *
06D6 86 20 EDU *
06D8 86 20 EDU *
06DA 86 20 EDU *
06DC 86 20 EDU *
06DE 86 20 EDU *
06E0 86 20 EDU *
06E2 86 20 EDU *
06E4 86 20 EDU *
06E6 86 20 EDU *
06E8 86 20 EDU *
06EA 86 20 EDU *
06EC 86 20 EDU *
06EE 86 20 EDU *
06F0 86 20 EDU *
06F2 86 20 EDU *
06F4 86 20 EDU *
06F6 86 20 EDU *
06F8 86 20 EDU *
06FA 86 20 EDU *

06C0 97 27 STA A XTMP+1
06C2 DE 26 LDX XTMP
06C4 32 PUL A
06C6 A7 00 STA A O,X
06C8 A7 01 PUL A
06CA AF 04 STA A 1,X
06CC B6 06 LD A UNRAVA
06CE A7 02 STA A 2,X
06D0 B6 06 LD A UNRAVA+1
06D2 A7 03 STA A 3,X
06D4 35 TXS
06D6 DE 28 LDX XTMP2
06D8 39 RTS
* SET UP NEW STACK POINTER
* INITIALIZE NEW STACK FRAGMENT
* RETURN ADDR
* OLD SP
* ADDR OF UNRAVL ROUTINE
* SET NEW SP
* RESTORE X-REG
* AND RETURN
* NO MORE ROOM FOR STACK FRAGMENT
* WYEBYE!
* ADDR OF UNRAVL ROUTINE
* UNRAVL EDU *
* UNRAVL STACK FRAGMENT.
* IF END OF FRAGMENT STILL EQUALS ENDPTR, SIMPLY RESET ENDPTR,
* OTHERWISE ADD SPACE TO FREE LIST
* SAVE X-REG
* XTMP2
* RESTORE OLD SP
* END OF FRAG = ENDPTR?
* UNRAVLE ENDPTR
* YEP, JUST UPDATE ENDPTR
* (BY ADDING 256)
* RESTORE X-REG
* AND RETURN
* FREE ALL CELLS WITHIN FRAGMENT
* UNRAVL EDU *
* CLR *
* BSR BLKFR
* BRA UNRVRT
* RETURN
* BLKFR EDU *
* FREE BLOCK, ARG-REG SPECIFY TOP OF BLOCK
* AT RETURN, X-REG, ARG-REG, AND XTMP=A-REG POINT TO BOTTOM OF BLOCK
* XTMP+1 IS UNSPECIFIED ON ENTRY AND RETURN
* STA A XTMP
* SUB B #4
* STA B XTMP+1
* LDX XTMP
* BRA BLKFR
* STA A CUR,X
* STA B CUR+1,X
* LDA B COR,X
* SUB #4
* BCC BLKFR
* LDA B FREPTR+1
* STA B COR+1,X
* LDA B FREPTR
* STA B CUR,X
* LDA # XTMP+1
* STA B FREPTR+1
* STA A FREPTR
* STA A XTMP
* FREE CELLS UNTIL DOWN TO 256-BYTE BOUNDARY
* STA A XTMP
* STA B COR,X
* LDA B FREPTR+1
* STA B CUR+1,X
* LDA B FREPTR
* STA B CUR,X
* LDA # XTMP+1
* STA B FREPTR+1
* STA A FREPTR
* STA A XTMP
* FREE CELLS UNTIL DOWN TO 256-BYTE BOUNDARY
* STA A XTMP
* STA B COR,X
* LDA B FREPTR+1
* STA B CUR+1,X
* LDA B FREPTR
* STA B CUR,X
* LDA # XTMP+1
* STA B FREPTR+1
* STA A FREPTR
* STA A XTMP

```

071B EQU
071D DE 2A
071E 26 26
0721 96 2F
0723 26 1B
0725 8D 07 B4
0728 96 75
072A 81 14
072C 24 EF
072E 96 74
0730 26 EF
0732 96 2E
0734 4A
0735 91 2C
0737 23 1F
0739 97 2E
073B 96 2F
073D
073F 80 04
0741 97 2F
0743 DE 2E
0745 DF 3A
0747 20 08
0749
074F DF 3A
0749 EE 02
074B DF 2A
074D DE 3A
074F
074F 4F 00
0751 6F 01
0753 6F 02
0755 6F 03
0757 3F
075B CE 08 E5
075B 7E 01 36
075E
075E 7C 00 3F
0761
0761 DF 26
0763 8D 88
0765 96 24
0767 A7 00
0769 96 27
076B A7 01
076D 96 3E
076F A7 02
0771 96 3F
0773 A7 03
0775 DF 3E
0777 DE 24
0779 3F
077A
077A DE 3E
077B DE 3E
077C DE 3E
077D DE 3E
077E DE 3E
077F DE 3E
0780 DE 3E
0781 DE 3E
0782 DE 3E
0783 DE 3E
0784 DE 3E
0785 DE 3E
0786 DE 3E
0787 DE 3E
0788 DE 3E
0789 DE 3E
078A DE 3E
078B DE 3E
078C DE 3E
078D DE 3E
078E DE 3E
078F DE 3E
0790 DE 3E
0791 DE 3E
0792 DE 3E
0793 DE 3E
0794 DE 3E
0795 DE 3E
0796 DE 3E
0797 DE 3E
0798 DE 3E
0799 DE 3E
079A DE 3E
079B DE 3E
079C DE 3E
079D DE 3E
079E DE 3E
079F DE 3E
07A0 DE 3E
07A1 DE 3E
07A2 DE 3E
07A3 DE 3E
07A4 DE 3E
07A5 DE 3E
07A6 DE 3E
07A7 DE 3E
07A8 DE 3E
07A9 DE 3E
07AA DE 3E
07AB DE 3E
07AC DE 3E
07AD DE 3E
07AE DE 3E
07AF DE 3E
07B0 DE 3E
07B1 DE 3E
07B2 DE 3E
07B3 DE 3E
07B4 DE 3E
07B5 DE 3E
07B6 DE 3E
07B7 DE 3E
07B8 DE 3E
07B9 DE 3E
07BA DE 3E
07BB DE 3E
07BC DE 3E
07BD DE 3E
07BE DE 3E
07BF DE 3E
07C0 DE 3E
07C1 DE 3E
07C2 DE 3E
07C3 DE 3E
07C4 DE 3E
07C5 DE 3E
07C6 DE 3E
07C7 DE 3E
07C8 DE 3E
07C9 DE 3E
07CA DE 3E
07CB DE 3E
07CC DE 3E
07CD DE 3E
07CE DE 3E
07CF DE 3E
07D0 DE 3E
07D1 DE 3E
07D2 DE 3E
07D3 DE 3E
07D4 DE 3E
07D5 DE 3E
07D6 DE 3E
07D7 DE 3E
07D8 DE 3E
07D9 DE 3E
07DA DE 3E
07DB DE 3E
07DC DE 3E
07DD DE 3E
07DE DE 3E
07DF DE 3E
07E0 DE 3E
07E1 DE 3E
07E2 DE 3E
07E3 DE 3E
07E4 DE 3E
07E5 DE 3E
07E6 DE 3E
07E7 DE 3E
07E8 DE 3E
07E9 DE 3E
07EA DE 3E
07EB DE 3E
07EC DE 3E
07ED DE 3E
07EE DE 3E
07EF DE 3E
07F0 DE 3E
07F1 DE 3E
07F2 DE 3E
07F3 DE 3E
07F4 DE 3E
07F5 DE 3E
07F6 DE 3E
07F7 DE 3E
07F8 DE 3E
07F9 DE 3E
07FA DE 3E
07FB DE 3E
07FC DE 3E
07FD DE 3E
07FE DE 3E
07FF DE 3E

```

```

07EB 8B 02
07EA A7 03
07EC A6 01
07EE 8B 02
07F0 A7 01
07F2 08
07F3 08
07F4 08
07F5 08
07F6 20 CB
07F8
07F8 39
07F9
07F9 DE FE
07F9 DF 70
07F9 DE 24
07FE DF 72
0801 96 3F
0803 46
0804 DE 3E
0806 24 08
0808
0808 DF 28
080A DE 72
080C 80 32
080E 7A 00 29
0811 DE 28
0813
0813 27 06
0815
0815 80 11
0817 22 FC
0819 25 ED
0819
0819
0819 DE 30
0819 8D 21
081F 0E 4C
0821 8D 1D
0823 DE 34
0825 8D 19
0827 39
0828
0828 A6 00
082A 97 72
082C A6 01
082E 97 73
0830 8B 02
0832 A7 01
0834 A6 03
0836 8A 02
0838 A7 01
083A 46
083B EE 02
083D 09
083E 09
083F 39
0840
0840 DF 4E
0842 2E 48
0844
0844 39
0845
0845 09
0846 A6 03
0848
0848 8B 02
084A 97 71
084C A6 02
084E 97 70
0850 96 4E
0852 A7 02
0854 9A 6F
0856 BF 4E
0858 24 03
085A 7C 00 6F
085D
085D 8A 02
085F A7 03
0861
0861 DE 70
0863 27 DF
0865 8A 71
0867 56
0868 25 D8
086A A6 03
086C 85 02
086E 2A D8
0870 EA 01
0872 E7 03
0874 8A 6F
0876 CA 02
0878 E7 01
087A 97 6F
087C A6 02
087E EA 00
0880 E7 02
0882 DA 6E
0884 E7 00
088A 97 6E
088B
088B DE 4E
088A 2F 05
088C
088C 9A 6F
088E 46
088F 25 1E
0891 EA 01
0893 C5 02
0895 26 CA
0897 A6 00
0899 2F 10
089A
089A DE 3E
089B DE 3E
089C DE 3E
089D DE 3E
089E DE 3E
089F DE 3E
08A0 DE 3E
08A1 DE 3E
08A2 DE 3E
08A3 DE 3E
08A4 DE 3E
08A5 DE 3E
08A6 DE 3E
08A7 DE 3E
08A8 DE 3E
08A9 DE 3E
08AA DE 3E
08AB DE 3E
08AC DE 3E
08AD DE 3E
08AE DE 3E
08AF DE 3E
08B0 DE 3E
08B1 DE 3E
08B2 DE 3E
08B3 DE 3E
08B4 DE 3E
08B5 DE 3E
08B6 DE 3E
08B7 DE 3E
08B8 DE 3E
08B9 DE 3E
08BA DE 3E
08BB DE 3E
08BC DE 3E
08BD DE 3E
08BE DE 3E
08BF DE 3E
08C0 DE 3E
08C1 DE 3E
08C2 DE 3E
08C3 DE 3E
08C4 DE 3E
08C5 DE 3E
08C6 DE 3E
08C7 DE 3E
08C8 DE 3E
08C9 DE 3E
08CA DE 3E
08CB DE 3E
08CC DE 3E
08CD DE 3E
08CE DE 3E
08CF DE 3E
08D0 DE 3E
08D1 DE 3E
08D2 DE 3E
08D3 DE 3E
08D4 DE 3E
08D5 DE 3E
08D6 DE 3E
08D7 DE 3E
08D8 DE 3E
08D9 DE 3E
08DA DE 3E
08DB DE 3E
08DC DE 3E
08DD DE 3E
08DE DE 3E
08DF DE 3E
08E0 DE 3E
08E1 DE 3E
08E2 DE 3E
08E3 DE 3E
08E4 DE 3E
08E5 DE 3E
08E6 DE 3E
08E7 DE 3E
08E8 DE 3E
08E9 DE 3E
08EA DE 3E
08EB DE 3E
08EC DE 3E
08ED DE 3E
08EE DE 3E
08EF DE 3E
08F0 DE 3E
08F1 DE 3E
08F2 DE 3E
08F3 DE 3E
08F4 DE 3E
08F5 DE 3E
08F6 DE 3E
08F7 DE 3E
08F8 DE 3E
08F9 DE 3E
08FA DE 3E
08FB DE 3E
08FC DE 3E
08FD DE 3E
08FE DE 3E
08FF DE 3E

```

0978 07 6F STA B SONPTR+1 NOPE, TRACE DOWN CAR
 0979 07 6E STA A SONPTR
 0980 07 6E LDA A DADPTR
 0981 A7 00 STA A CAR, X
 0982 07 6E LDA A DADPTR+1
 0983 08 02 ORA A MRRKBIT
 0984 07 01 STA A CAR+1, X
 0985 20 26 BRA GCOC44

JOIN COMMON CODE

GCOC44 EQU *
 * CAR IS H/L OR NUM, TRACE DOWN CDR
 LDA A CDR+1, X
 BRA GCOC43 (JUST LIKE AN ATOM CELL)

GCOC44 EQU *
 * FORM IS ATOM
 LDA A CDR+1, X ATOM ALREADY MARKED?
 BIT A MRRKBIT
 RNE GCOC44
 DEX
 LDA B CAR+1, X
 COMPLEMENT GC BIT OF CAR
 GCOC43 EQU *
 * AT THIS POINT C-BIT SET IF WE HAVE AN ATOM...
 EOR B MRRKBIT
 STA B CAR+1, X
 LDA B CDR, X
 BLE GCOC42
 BLE GCOC42
 STA B SONPTR+1
 LDA A DADPTR
 STA A CDR, X
 LDA A DADPTR+1
 ORA A MRRKBIT
 STA A CDR+1, X
 RCC GCOC44

GCOC44 EQU *
 * DO WE HAVE AN ATOM?
 YEP, SET LOW BIT AGAIN
 STX
 BRA GCOC42 JOIN COMMON CODE

GCOC44 EQU *
 * STACK UNDERFLOW
 FCB 4
 MBOCNG FCC 4
 *NO MORE FREE SPACE
 FCB 4
 STKENG FCC 4
 *NO ROOM FOR STACK
 FCB 4

REGADR EQU *
 END

NO ERROR(S) DETECTED

SYMBOL TABLE:

ALP	0044	ATMINI	0139	REGADR	090A	REGPTR	0020	BLKF1	0704
BLKF2	070A	BLKFE	04FA	CAR	0000	CDR	0002	CELPTR	003A
CHP12	0042	CHPTR	0040	CONATH	0060	CUREVL	006C	DADPTR	0070
DADPTR	0050	EDHMR	00FC	EMDPR	002E	EVAL	0004	EDATH	0084
EDSTKM	00D5	ERRRKR	0154	ERREX	0136	FORM	0100	EVLATH	0118
FLP	004B	FORM	0034	FRECL	07A9	FREPTR	002A	FRNUM	014B
GF1	07C3	GF2	07D2	GFCLR	07E8	GFNUM	07F8	GFNUM	07F2
GFREE	0074	GCL2	0730	GCOC2	08AF	GCOC2	0888	GCOC3	0888
GCOC4	08D1	GCOC4	088F	GCOC2	085D	GCOCDD	0848	GCODAT	0845
GCODUN	0861	GCOC6	088C	GCOL	0784	GCOT1	0809	GCOT2	0813
GCOT3	0815	GCOT4	0818	GCOTEMP	0072	GCOTHR	0828	GCOTRJ	0840
GCOTRCE	07F9	GCOTR	0844	GETC	0142	GETCAR	0163	GETCEL	0710
GETSYM	0157	GFNXL	0127	GNYXL	0124	ISATOM	013F	ISDTPR	013C
ISVAR	079A	ISVND	0767	LAMATH	005A	LBRKAT	0068	LOOKUP	0121
LPARAT	004C	LSTAD2	0130	LSTADD	0120	LSTENO	077A	LSTEND	0133
LSTINI	012A	LSTPTR	0036	HINFE	0014	MRRBIT	0002	MANPTR	0032
MEGET	07AF	MLATH	004A	MLAMAT	0058	NLP	0046	NOGCEL	0758
NOGNG	08E3	MSUBAT	005C	MUMFRM	0148	MUMINV	0145	OREGAD	06A0
OLGET	0747	PEEK	00F0	POPX	077A	POPXR	07BD	FRINR	0149
PRINT	0108	PROPOP	077A	PROPSH	075E	PUSHX	0781	PUTC	0151
POTSVM	014E	QUATN	0054	RKAT	006A	READ	0103	APARAT	004E
RSETAT	0042	RSIFLG	00F1	SETATM	011E	SONPTR	006E	SFCTPR	002C
SFSAVE	003C	SOUTAT	0052	STKENG	08F8	STKF2	0A49	STKF4	0685
STKFR6	06A0	STKFUL	040A	STKIM	0032	STKPTR	003E	STPF1	002A
STP2	002B	STP1	0022	STP2	002A	SUBLS2	0066	SUBLS3	002A
SUBRAT	005A	SYHLST	0030	SYMPTR	003B	TATOH	005E	TOPX	0793
UNRAVA	06E0	UNRAVL	06E2	UNRFRE	06F5	UNRVRT	06F2	XIMP2	0028
XIMP2	0028	ZERO	00FE						

0940 0940 00 F1
 0941 00 0C
 0942 26 0C
 0943 00 00
 0944 2F 02
 0945 00 00
 0946 0F 34
 0947 00 00
 0948 0F 34
 0949 00 00
 0950 39
 0951 0E 64
 0952 20 F9
 0953 00 00
 0954 DF 34
 0955 DF 34
 0956 00 00
 0957 00 00
 0958 00 00
 0959 00 00
 0960 00 00
 0961 00 00
 0962 00 00
 0963 00 00
 0964 00 00
 0965 00 00
 0966 00 00
 0967 00 00
 0968 00 00
 0969 00 00
 0970 00 00
 0971 00 00
 0972 00 00
 0973 00 00
 0974 00 00
 0975 00 00
 0976 00 00
 0977 00 00
 0978 00 00
 0979 00 00
 0980 00 00
 0981 00 00
 0982 00 00
 0983 00 00
 0984 00 00
 0985 00 00
 0986 00 00
 0987 00 00
 0988 00 00
 0989 00 00
 0990 00 00
 0991 00 00
 0992 00 00
 0993 00 00
 0994 00 00
 0995 00 00
 0996 00 00
 0997 00 00
 0998 00 00
 0999 00 00

REGADR EQU *
 END

NOPE, TRACE DOWN CAR
 THIS SHOULD BE IN EACH MODULE
 * DIRECT PAGE STORAGE CELLS

ENSURE THAT FORN STARTS OUT AS LEGAL FORM

NO ERROR(S) DETECTED

SYMBOL TABLE:

ALP	0044	ATMINI	0139	REGADR	090A	REGPTR	0020	BLKF1	0704
BLKF2	070A	BLKFE	04FA	CAR	0000	CDR	0002	CELPTR	003A
CHP12	0042	CHPTR	0040	CONATH	0060	CUREVL	006C	DADPTR	0070
DADPTR	0050	EDHMR	00FC	EMDPR	002E	EVAL	0004	EDATH	0084
EDSTKM	00D5	ERRRKR	0154	ERREX	0136	FORM	0100	EVLATH	0118
FLP	004B	FORM	0034	FRECL	07A9	FREPTR	002A	FRNUM	014B
GF1	07C3	GF2	07D2	GFCLR	07E8	GFNUM	07F8	GFNUM	07F2
GFREE	0074	GCL2	0730	GCOC2	08AF	GCOC2	0888	GCOC3	0888
GCOC4	08D1	GCOC4	088F	GCOC2	085D	GCOCDD	0848	GCODAT	0845
GCODUN	0861	GCOC6	088C	GCOL	0784	GCOT1	0809	GCOT2	0813
GCOT3	0815	GCOT4	0818	GCOTEMP	0072	GCOTHR	0828	GCOTRJ	0840
GCOTRCE	07F9	GCOTR	0844	GETC	0142	GETCAR	0163	GETCEL	0710
GETSYM	0157	GFNXL	0127	GNYXL	0124	ISATOM	013F	ISDTPR	013C
ISVAR	079A	ISVND	0767	LAMATH	005A	LBRKAT	0068	LOOKUP	0121
LPARAT	004C	LSTAD2	0130	LSTADD	0120	LSTENO	077A	LSTEND	0133
LSTINI	012A	LSTPTR	0036	HINFE	0014	MRRBIT	0002	MANPTR	0032
MEGET	07AF	MLATH	004A	MLAMAT	0058	NLP	0046	NOGCEL	0758
NOGNG	08E3	MSUBAT	005C	MUMFRM	0148	MUMINV	0145	OREGAD	06A0
OLGET	0747	PEEK	00F0	POPX	077A	POPXR	07BD	FRINR	0149
PRINT	0108	PROPOP	077A	PROPSH	075E	PUSHX	0781	PUTC	0151
POTSVM	014E	QUATN	0054	RKAT	006A	READ	0103	APARAT	004E
RSETAT	0042	RSIFLG	00F1	SETATM	011E	SONPTR	006E	SFCTPR	002C
SFSAVE	003C	SOUTAT	0052	STKENG	08F8	STKF2	0A49	STKF4	0685
STKFR6	06A0	STKFUL	040A	STKIM	0032	STKPTR	003E	STPF1	002A
STP2	002B	STP1	0022	STP2	002A	SUBLS2	0066	SUBLS3	002A
SUBRAT	005A	SYHLST	0030	SYMPTR	003B	TATOH	005E	TOPX	0793
UNRAVA	06E0	UNRAVL	06E2	UNRFRE	06F5	UNRVRT	06F2	XIMP2	0028
XIMP2	0028	ZERO	00FE						

```

096F 96 34 LDA A FORM
0971 A7 00 STA A CAR,X
0973 96 35 LDA A FORM+1
0975 A7 01 STA A CAR+1,X
0977 96 6C LDA A CUREVL
0979 A7 02 STA A CDR,X
097B 96 6D LDA A CUREVL+1
097D A7 03 STA A CDR+1,X
097F DF 6C STX CUREVL
0981 DF 34 LDX FORM
0983 EE 02 CDR,X
0985 DF 44 STX ALP
0987 DF 34 LDX FORM
0989 EE 00 CAR,X
098B DF 83 BSR EVAL
098D DF 0A E3 JSR ISDIRR
098F 25 1A BCS EVLERR
* NOW X-REG SHOULD BE EXPLICIT LAMBDA EXPR
0992 DF 34 STX FORM
0994 EE 02 CDR,X
0996 DF 44 STX FLP
0998 DF 34 LDX FORM
099A EE 00 CAR,X
099C 9C 5A CPX SURRAT
099E 27 18 BCD ELSURAT
09A0 9C 5C CPX NSURAT
09A2 27 0D BCD ELSUR
09A4 9C 56 CPX LAHATH
09A6 27 6D BCD ELLAM
09A8 9C 58 CPX MLAMAT
09AA 27 5C BCD EVLMA
09AC EVLERR EQU *
09AE CE 0B BC LDX NEVLMS
09AF 20 33 BRA EVLSR2
*
09B1 EVLSUR EQU *
* EVALUATE (NSUR,FUNCAORR)
09B2 DE FE LDX ZERO
09B3 DF 01 60 JSR PROPSH
09B6 20 09 BRA EVNSU2
*
09B8 EVLSUR EQU *
* EVALUATE (SUBR,FUNCAORR)
09BB DF 2F BSR EVLALS
09BA DE 64 LDX EDIATH
09BC 7D 00 F1 TST RSTFLG
09BE 26 18 BNE EVLSR5
09C1 EQU *
09C1 DE 48 LDX FLP
09C3 8D 0A EE JSR TSATON
09C6 25 17 BCS EVLSER
09C8 2F 15 BLE EVLSER
09CA 09 DEX
09CB EE 00 LDX CAR,X
09CD AD 00 LIA A 0,X
09CF 81 21 CMP A #1
09D1 26 0C BNE EVLSER
09D3 4D 01 TST 1,X
09D5 26 08 BNE EVLSER
* ALP NOW HAS ARG LIST
09D7 AD 02 JSR 2,X
09D9 EVLSRS EQU *
09DB DF 34 STX FORM
09DB DF 69 BSR POPFRE
09DD 2D 4A BRA EVLRT5
09DF EVLSER EQU *
09DF DF 65 BSR POPFRE
09E1 CE 0B 6E LDX NEVLMS
09E4 EVLER2 EQU *

```

```

09E4 DF 01 54 JSR ERRBRK
09E7 20 3E BRA EVLRT4
*
09E9 EVLALS EQU *
09E9 DF 0A FA JSR LSTINI
09EC DE 44 LDX ALP
09EE DF 0R 33 JSR GNXTL
09F1 25 0C BCS EVLSR3
09F3 DF 44 STX ALP
09F5 DE 26 LDX XTMP
09F7 DF 09 40 JSR EVAL
09FA DF 0B 06 BSR LSTARD
09FC 20 ED BRA EVLSR2
09FF EQU *
09FF DF 01 15 LSTENO
0A02 DF 01 1B JSR TOPX
0A05 DF 44 STX ALP
0A07 39 RTS
*
0A08 EVLMA EQU *
* EVALUATE (LAMBDA (L) EXPR)
* SIMULATE EVALS BY MAKING ONE-ELEMENT LIST
* WITH UNEVAL'D ARG LIST AS SINGLE "ACTUAL" ARG.
0A0B DF 0A FA JSR LSTINI
0A0D DE 44 LDX ALP
0A0F DF 0B 04 JSR LSTARD
0A10 DF 01 15 LDX LSTENO
0A13 20 02 BRA EVLN2
* JOIN COMMON CODE
0A15 EVLLAM EQU *
* EVALUATE (LAMBDA (X Y) EXPR)
0A15 DF 02 BSR EVLALS
0A17 EQU *
0A17 DF 6A BSR EVLSV
* TOP OF STACK NOW POINTS TO LIST OF FORMALS
* JUST BELOW IS LIST OF THEIR OLD VALUES
0A19 EVLROD EQU *
0A19 DE 48 LDX FLP
0A1B DF 01 63 JSR GETCAR
0A1E DF 09 40 JSR EVAL
0A21 DF 48 STX FLP
0A23 DF 38 BSR EVLRST
0A25 DE 48 LDX FLP
0A27 DF 34 EQU *
0A27 DF 34 STX FORM
0A29 EQU *
0A29 DE 6C LDX CUREVL
0A2B DF 25 BSR GFNXTA
0A2D DE 44 LDX ALP
0A2F DF 6C STX CUREVL
0A31 32 EQU *
0A31 32 PUL A
0A32 97 48 STA A FLP
0A34 32 EQU *
0A34 32 PUL A
0A35 97 49 STA A FLP+1
0A37 32 EQU *
0A37 32 PUL A
0A38 97 46 STA A MLP
0A3A 32 EQU *
0A3A 32 PUL A
0A3B 97 47 STA A HLF+1
0A3D 32 EQU *
0A3D 32 PUL A
0A3E 97 44 STA A ALP
0A40 32 EQU *
0A40 32 PUL A
0A41 97 45 STA A ALP+1
0A43 DE 34 LDX FORM
0A45 39 RTS
*
0A46 POPFRE EQU *
* FREE THE ARG LIST (LIST HEAD IS AT TOP OF STACK)
0A46 DF 01 6C JSR PROPOP
*
0A48 DF 44 ADVANCE ALP
0A49 DF 00 LDX CAR,X
0A50 DF 28 STX XTMP2
* GET OLD VALUE OF FORMAL ARG, SAVE IN LIST
0A52 DE 26 LDX XTMP
0A53 DF 01 66 JSR ISVAR
0A54 25 1B BCS EVLNER
0A56 DF 09 4C JSR EVALM
* OLD VALUE NOW IN FORM, SAVE IN LIST
0A58 DE 44 LDX ALP
0A59 96 34 LIA A FORM
0A5A 47 00 STA A CAR,X
0A5B 96 35 LIA A FORM+1
0A5C A7 01 STA A CAR+1,X
* SET UP NEW VALUE FOR ATON
0A5E DE 28 LDX XTMP2
0A5F DF 34 STX FORM
0A60 DF 26 LDX XTMP
0A61 DF 0B 65 JSR SETATM
0A62 DE 46 LDX NLP
0A63 24 8D BNE EVLNS1
0A65 39 RTS
* WE GOT A BAD FORMAL ARG, COMPLAIN...
0A66 DF 34 JSR FORM
0A67 CE 0B 7D LDX NEVLMS
0A68 DF 01 54 JSR ERRBRK
0A6E 20 F5 BRA EVLNS5
*
0A6E EQU *
0A6E DF 26 SEC
0A6F DF 26 STX XTMP
0A70 2F 05 BLE ISDIRT
0A71 96 27 LIA A XTMP+1
0A72 46 ROR A
0A73 DE 26 LDX XTMP
0A74 39 RTS
* RETURN C-BIT SET IF NOT DTPR
* RETURN Z-BIT SET IF IS NIL
* BCC WILL BRANCH ON DTPR
* BEO WILL BRANCH ON NIL
0A75 00 EQU *
0A75 00 STX XTMP
0A76 2F 05 BLE ISDIRT
0A77 96 27 LIA A XTMP+1
0A78 46 INC A
0A79 46 ROR A
0A7A DE 26 LDX XTMP
0A7B 39 RTS
* ISATON EQU *
* RETURN WITH C-BIT IF NOT ATON
* RETURN WITH Z-BIT IF NIL
* RETURN WITH N-BIT SET IF NUMBER
* BHI WILL BRANCH ON NON-NIL ATON
0A7C 00 EQU *
0A7C 00 STX XTMP
0A7D 2F 06 BLE ISDIRT
0A7E 96 27 LIA A XTMP+1
0A7F 46 INC A
0A80 46 ROR A
0A81 DE 26 LDX XTMP
0A82 39 RTS
*
0A82 EQU *
0A82 DF 01 0C JSR GETCEL
0A83 DE 44 LDX ALP
0A84 96 3A LIA A CELFK
0A85 47 02 STA A CAR,X
0A86 96 3B LIA A CELPK+1
0A87 A7 03 STA A CAR+1,X
0A88 DF 34 LDX CELPKR
0A89 EQU *

```



```

0801 09 DEX
0802 09 DEX
0803 7E 01 12 JMP PUSH FUSH POINTER AND RETURN
0806
0806 DF 34 LSTAD2 EQU *
* ADD FORM IN X-REG TO LIST POINTER ON STACK
* CALL LSTAD2 IF ALREADY STORED IN 'FORM'
0808 8D 01 0C LSTAD2 EQU *
0808 8D 01 0C JSR GETCEL GET A NEW CELL
0808 96 34 LDA A FORM FILL IN CAR
0809 A7 00 STA A CAR,X
0809 A7 00 STA A CAR,X
0809 96 35 LDA A FORM+1
0811 A7 01 STA A CAR+1,X
0813 8D 01 18 JSR TOPX GET LIST POINTER
0816 96 34 LDA A CELPTR (CELPTR FILLED IN BY GETCEL)
0818 A7 02 STA A CAR,X
0818 A7 02 STA A CAR,X
081C E7 03 STA B CAR+1,X
081E DE 3E LDX STKPTR
0820 A7 00 STA A CAR,X
0822 E7 01 STA B CAR+1,X
0824 39 RTS
0825
0825 DF 34 LSTEN2 EQU *
* FILL IN CAR OF LAST CELL IN LIST
* ALSO POP OFF LIST POINTER
* CALL LSTEN2 IF FORM ALREADY STORED IN 'FORM'
0827
0827 8D 01 15 LSTEN2 EQU *
0828 96 34 JSR POPX POP OFF LIST POINTER
082C A7 02 STA A FORM
082E 96 35 LDA A FORM+1
0830 A7 03 STA A CAR+1,X
0832 39 RTS
0115 LSTENO EQU POPX USE THIS IF LIST ENDED WITH NIL
0833
0833 DF 26 GNTXL EQU *
* GET NEXT LIST ELEMENT
* X-REG CONTAINS LIST POINTER
* RETURNS CAR IN X-REG, CAR IN XTMP, C-BIT SET IF NOT DIFR
0835 2F 10 BLE GNLNO
0837 96 27 LDA A XTMP+1
0839 46 ROR A
083A 25 08 RCS GNLNO
083C A6 00 LDA A CAR,X
083E 97 26 STA A XTMP
0840 A6 01 LDA A CAR+1,X
0842 97 27 STA A XTMP+1
0844 EE 02 LDX CAR,X
0846 39 RTS
0847 EQU *
0847 DE FE GNLNO
0849 DF 26 LDX ZERO
084B 8D SEC
084C 39 RTS
084D
084D 8D 0A E3 GFNXTL EQU *
* SAME AS GNTXL, EXCEPT ALSO FREE CELL AFTER
* RETURNING ITS CONTENTS
0850 25 F5 JSR ISDTPR
0852 A6 02 RCS GNLNO
0854 97 28 LDA A CAR,X
0856 A6 03 STA A XTMP+1
0858 97 29 STA A XTMP2+1

```

NO ERROR(S) DETECTED

SYMBOL TABLE

```

ALP 0044 ATHINI 0139 BEGADR 08A2 BEOPTR 0020 CAR 0000
CAR 0002 CELPTR 003A CMPTM2 0042 CMPTM 0040 CONATH 0060
CUREVL 004C DADPTR 0070 ENDMEN 00FC ENDFTR 002E
EOL 0004 EOIATM 0064 ERBRK 0154 ERREX 0136 EVAL 0940
EVALS 09E9 EULATH 094C EULDR 0A19 EULDTF 0955 EULENS 08BC
EULEI 0951 EVLER2 09E4 EVLERR 09AC EULLAM 0A15 EULNER 0A09
EVLN2 0A17 EVLNLA 0A09 EVLNMS 0B7B EVLMS1 0A95 EULMS2 0A81
EVLNS 0A08 EVLNSU 0981 EVLNSV 0A83 EVLRSO 0A62 EULRT4 0A27
EVLRS2 0A49 EVLRS3 0A6D EVLRSI 095B EVLRT 094E EULRT4 0A27
EVLRS 0A29 EULSR2 09FC EULSR3 09FF EULSR5 09D9 EULSER 07DF
EVLMS 08AE EVLSOB 098B EWSU2 09C1 FLP 0048 FORM 0034
FRECEL 010F FREPTR 002A FRMHU 014B GCFREE 0074 GCOL 015D
GCTEMP 0072 GETC 0142 GETCAR 0163 GETSTH 0157
GFNXTA 0A52 GFNXTL 084D GNLNO 0847 GNTL 0833 ISATOM 0AEE
ISAIPT 0A99 ISDTPR 0A63 ISDTR 0AED ISVAR 0166 LAMATH 0054
LBRKAT 006B LOOKUP 0121 LPARAT 004C LSTAD2 080B LSTAD 080A
LSTENO 0115 LSTEN2 0827 LSTEND 0825 LSTINI 0AFA LSTINX 0801
LSTPTR 0036 MANPTR 0032 ATLATA 004A MLAMAT 0038 MLP 0046
NSUBAT 005C MUMFRM 0148 NUMINV 0145 OBEGAD 0940 PEEKC 00F0
PFR2 0A4B PFR3 0A4D PFERUN 0A51 POPFRE 0A46 POPX 0115
PRMR 0169 PRINT 010A PROPOP 01AC PROFSH 0160 PUSH 0112
PUIC 0131 PUTSYM 014E QUATH 0054 RBRKAT 004A READ 0103
RPARAT 004E RSETAT 0062 RSTFLG 00F1 RSETAT 0865 SMDPTR 006E
SMDPTR 002C SPSAVE 003C SOUTAT 0052 STKFRG 015A STKPTR 003E
SIF1 0026 SIF2 0028 STR1 0022 STR2 0024 SUBLS2 0064
SUBLS3 0026 SURRAT 005A SYMLST 0030 SYMPTR 003B TATDM 005E
TOPX 0118 XTMP 0026 ZERO 00FE

```

```

0860 UREGAD EQU #980
0800 CAR EQU 0
0802 CDR EQU 2
0804 EOI EQU 4
* DIRECT PAGE STORAGE CELLS
0820 ORG 420
0820 0E 7D BEOPTR FDB BEGADR THIS SHOULD BE IN EACH MODULE
0822 STK1 RMB 2
0824 STK2 RMB 2
0826 STK3 RMB 2
0828 STK4 RMB 2
082A STK5 RMB 2
082C STK6 RMB 2
082E STK7 RMB 2
0830 STK8 RMB 2
0832 STK9 RMB 2
0834 STK10 RMB 2
0836 STK11 RMB 2
0838 STK12 RMB 2
083A STK13 RMB 2
083C STK14 RMB 2
083E STK15 RMB 2
0840 STK16 RMB 2
0842 STK17 RMB 2
0844 STK18 RMB 2
0846 STK19 RMB 2
0848 STK20 RMB 2
084A STK21 RMB 2
084C STK22 RMB 2
084E STK23 RMB 2
0850 STK24 RMB 2
0852 STK25 RMB 2
0854 STK26 RMB 2
0856 STK27 RMB 2
0858 STK28 RMB 2
0860 STK29 RMB 2
0862 STK30 RMB 2
0864 STK31 RMB 2
0866 STK32 RMB 2
0868 STK33 RMB 2
086A STK34 RMB 2
086C STK35 RMB 2
086E STK36 RMB 2
0870 STK37 RMB 2
0872 STK38 RMB 2
0874 STK39 RMB 2
0876 STK40 RMB 2
* SINGLE CHAR SLOTS ...
08F0 ORG #FO
08F0 PEEKC RMB 1
08F1 RSTFLG RMB 1
* GLOBAL CONSTANTS
08FC ENHLEN RMB 2
08FE ZERO FDB 0
* GLOBAL VECTORS
0100 ORG #100
0100 EVAL RMB 3
0103 READ RMB 3
0106 PRINT RMB 3
0109 GETCEL RMB 3
010C FRECEL RMB 3
010F

```

```

0112 PUSHX RMB 3
0115 POPX RMB 3
0118 TOPX RMB 3
011B EVLAIN RMB 3
011E SETAIN RMB 3
0121 LOOKUP RMB 3
0124 GNTXL RMB 3
0127 GFAXTL RMB 3
012A LSTINI RMB 3
012D LSTAD2 RMB 3
0130 LSTAD RMB 3
0133 LSTEND RMB 3
0135 LSTENO EQU POPX
0136 ERKEX RMB 3
0139 JHP ATHINI
013C ISDTPR RMB 3
013F ISATOM RMB 3
0142 GETC RMB 3
0145 NUMINV RMB 3
0148 FRMHU RMB 3
014E PUTSYM RMB 3
0151 ERBRK RMB 3
0154 BEYSTA RMB 3
0157 STKFRG RMB 3
015A SCOL RMB 3
015D PROFSH RMB 3
0160 GETCAR RMB 3
0163 ISVAR RMB 3
0166 FRMR RMB 3
0169 PROPOP RMB 3
0880 ORG 0860
0880 UARMS EQU #7103
0880 NOTHAT EQU FLP
*
0880 ATHINI EQU *
* INITIALIZE BUILT-IN ATOMS
* IN THE FORM OF NULL-TERM'D STRING
* FOLLOWED BY 8-BIT ADDRESS
0880 CE 00 88 LDA #ATLST POINT TO BEGINNING OF ATOMS
0883 DF 38 STX STKPTR
0885 EQU *
0885 8D 01 21 JSR LOOKUP LOOK UP, AND PUT ON SYMLST
* (RESULT FROM LOOKUP ALSO IN FORM)
0888 DE 38 LDX SYMPTR SCAN TO END OF ATOM NAME
088A 08 EQU *
088A 08 INX
088B 40 00 TST 0,X
088D 26 FB BNE AT13
088F DF 38 STX SYMPTR
0891 EE 00 LDX 0,X
0893 96 34 LDA A FORM
0895 A7 00 STA A 0,X
0897 96 35 LDA A FORM+1
0899 A7 01 STA A 1,X
089B DE 38 LDX SYMPTR
089D 08 INX
089E 08 INX
089F 40 00 TST 0,X
08A1 26 E2 BNE AT12
* YEP, ALL HOME
* INITIALIZE EOIATH
08A3 CE 00 E0 LDX #EOIATH
08A6 8D 01 21 JSR LOOKUP
08A9 DF 84 STX EOIATH
* INITIALIZE "NOTHING" AND EOIATH TO POINT TO EOIATH
08AB 8D 28 JSR #SETAT

```

0BDD DE 48 LDX NOTHAT
 0BDF BD 27 RSR SSETAT
 0BE1 DE 5C LDX TATOM
 0BE3 EF 01 STX CRR-1,X
 0BES CE 0D E2 LDX #SUBRLS INITIALIZE PRIMARY SUBRS AND #SUBRS
 0BE8 BD 21 RSR SXSINI
 0BEA DE 5A LDX SUBRAT COPY SELFUM TO LAMBDA/LAMBDA, ETC.
 0BEC EE 01 LDX CRR-1,X
 0BEF DF 34 STX FORN
 0BF0 DE 5C LDX #SUBRAT
 0BF2 BD 14 RSR SSETAT
 0BF4 DE 5A LDX LAMATH
 0BF4 BD 10 RSR SSETAT
 0BF8 DE 5B LDX NLANAT
 0BFA BD 0C RSR SSETAT
 0BFC DE 40 LDX COMATH
 0BFE BD 0B RSR SSETAT
 0C00 DE 66 LDX SUBLS2
 0C02 BD 07 RSR SXSINI
 0C04 DE 76 LDX SUBLS3
 0C06 20 03 BRA SXSINI
 0C0B EQU SSETAT EQU * SETATH ALLOW FOR RSR'S
 0C0B 7E 01 IE *
 0C0B EQU SXSINI EQU *
 0C0B CA 5A LDX B #SUBRAT INITIALIZE SUBRS
 0C0D BD 03 RSR SUBRINI
 0C0F 0B INX INX POINT TO NSUBR LIST
 0C10 CA 5C LDX B #NSUBRAT AND RETURN THROUGH SUBINI
 0C12 SUBRINI EQU *
 0C12 EQU * B-REG HOLDS ADDR OF ATOM (IN DIR. PAGE)
 0C12 EQU * K-REG POINTS TO LIST OF NAMES
 0C12 EQU * FOLLOWED BY ADDRESS OF FUNCTION
 0C12 EQU * FUNCTION MUST START WITH FCB 'I'; FCB 0
 0C12 STX SYMPTR
 0C14 7F 00 24 CLR XTMP
 0C17 D7 27 STA B XTMP+1
 0C19 DE 26 LDX XTMP
 0C1B EE 00 LDX 0,X
 0C1D DF 44 STX ALP
 0C1F EQU SBLUP EQU *
 0C1F DE 3B LDX SYMPTR POINT TO ATOM NAME
 0C21 6D 00 TST 0,X ALL DONE?
 0C23 27 4D BEG SBLUP YUP
 0C25 BD 01 21 JSR LOOKUP FORM ATOM (RETURNED IN 'FORM')
 0C28 DF 46 STX NLP SAVE IT FOR LATER
 0C2A DE 3B LDX SYMPTR SKIP PAST ATOM NAME
 0C2C EQU SBLUP2 EQU *
 0C2C 0B INX
 0C2D 6D 00 TST 0,X
 0C2F 26 FB RNE SBLUP2
 0C31 0B INX
 0C32 DF 3B STX SYMPTR
 0C34 EE 00 LDX 0,X
 0C36 A6 00 LDX A 0,X
 0C38 81 21 CE3B 81 21 CNP A B 'I'
 0C3A 26 37 RNE SBLRAD
 0C3C 6D 01 TST 1,X
 0C3E 26 33 RNE SBLRAD
 0C40 DF 4B STX FLP
 0C42 BD 01 0C JSR GETCEL
 0C45 9A 4B LDX A FLP
 0C47 A7 00 STA A CAR.X

0C49 96 49 LDX A FLP+1
 0C4B A7 01 STA A CAR+1,X
 0C4D 0B INX
 0C4E DF 34 STX FORN
 0C50 BD 01 0C JSR GETCEL
 0C53 96 44 LDX A ALP
 0C55 A7 00 STA A CAR,X
 0C57 96 45 LDX A ALP+1
 0C59 A7 01 STA A CAR+1,X
 0C5B 96 34 LDX A FORN
 0C5D A7 02 STA A CRR,X
 0C5F 96 35 LDX A FORM+1
 0C61 A7 03 STA A CRR+1,X
 0C63 DF 34 STX FORN
 0C65 DE 46 LDX NLP
 0C67 BD 01 IE JSR SETATH
 0C6A DE 3B LDX SYMPTR
 0C6C 0B INX
 0C6E 0B INX
 0C6F DF 3B STX SYMPTR
 0C70 20 AB BRA SBLUP
 0C72 EQU SPIRT EQU *
 0C72 39 EQU *
 0C73 EQU SBLRAD EQU *
 0C73 DE 0E 3C LDX #SBIENS
 0C76 7E 01 36 JMP ERREX
 2100 EQU * BUILT-IN FUNCTIONS
 0C79 EQU MAGUD EQU * \$2100
 0C79 EQU SELFUN EQU *
 0C79 21 00 EQU * RETURN SELF AS VALUE (USED FOR LAMBDA/NLAMBDA/SUBR/NSUBR/COMT)
 0C7B DE 6C LDX CUREVL
 0C7B EE 00 LDX CAR,X
 0C7F 39 EQU RTS
 0C80 EQU SYSFUN EQU *
 0C80 21 00 EQU * (SYS) RETURN TO RUS
 0C82 7E 71 03 JMP WARMs
 0C85 EQU EVLFUN EQU *
 0C85 21 00 EQU * GET FIRST ARG
 0C87 BD 74 BSR GCARA AND EVAL IT, AND RETURN
 0C89 7E 01 00 JMP EVAL
 0C8C EQU REPFUN EQU *
 0C8C 21 00 EQU *
 0C8E 7E 01 03 JMP READ
 0C91 EQU PRIFUN EQU *
 0C91 21 00 EQU * GET FIRST ARG
 0C93 BD 68 BSR GCARA
 0C95 7E 01 06 JMP PRINT
 0C9B EQU CARFUN EQU *
 0C9B 21 00 EQU * GET FIRST ARG
 0C9A BD 61 RSR GCARA
 0C9C BD 01 3C JSR ISDTPR
 0C9F 27 05 BEQ CARNIL
 0CA1 25 12 RCS CAREER
 0CA3 EE 00 LDX CAR,X
 0CA5 39 EQU RTS
 0CAA EQU CARMIL EQU *
 0CA6 DE 64 LDX EDIATH
 0CAB 39 EQU RTS
 0CAB EQU CONFUN EQU *
 0CAB 21 00 EQU * GET FIRST ARG
 0CAB BD 50 RSR GCARA

0CAB 8D 50 RSR GCARA

0F8A DE 2C	LDX SPCPTR	SAVE POINTER TO BEGINNING OF NEW NAME
0F8C DF 3B	STX SYMPTR	
0F8E FE 10 22	LDX GNNLAS	GET LAST ATOM NAME USED
0FC1 24 05	RNE GNM2	
0FC3 CE 10 24	LDX MEMFIR	NEVER INITIALIZED, USE FIRST
0FC6 20 03	BRA GNM3	
0FCB 09	EDU *	
0FC9 EE 00	LDX CAR,X	POINT TO CELL
0FCB	EDU *	POINT TO NAME
0FCB A6 00	LDX TO NAME SPACE	
0FCD DF 2A	LDX STX	
0FCE DE 2C	LDX SPCPTR	
0FB1 9C 2E	CPX ENDPTR	
0FB3 2A 47	BPL GNMHAD	
0FB5 A7 00	STA A 0,X	
0FB7 0B	INX	
0FB8 DF 2C	STX SPCPTR	
0FBA 4D	TST A	ALL DONE?
0FDB 27 05	BEQ GNM4	
0FDD DE 2A	LDX XTMP	
0FDF 0B	INX	
0FE0 20 E9	BRA GNM3	NOPE, LOOP UNTIL NULL BYTE
0FE2 DE 3B	LDX EDU	POINT TO FIRST VARIABLE CHAR
0FE4 0B	INX	
0FE5 0B	INX	
0FE6 0B	INX	
0FE7	EDU *	
0FE7 A6 00	LDX A 0,X	
0FE9 2A 10	RNE GNM42	
0FE9 B6 41	* UP TO NULL BYTE, SET TO 'A'	
0FEA 0B	LDX A B'A	
0FE9 0B	STA A 0,X	
0FEF 0B	INX	
0FF0 9C 2E	CPX ENDPTR	
0FF2 2A 2B	BPL GNMHAD	
0FF4 AF 00	CLR	
0FF6 0B	INX	
0FF7 DF 2C	STX SPCPTR	
0FF9 20 0E	BRA GNM51	GO CHECK IF REALLY NEW
0FFB AC	EDU *	
0FFC B1 5A	INC A	
0FFE 23 07	CPX B'Z	UP TO Z?
1000 B6 41	RLS GNM5	
1002 A7 00	LDX A B'A	
1004 0B	INX	YES, RESET TO AA
1005 20 E0	BRA GNM41	GO ON TO NEXT BYTE
1007	EDU *	
1007 A7 00	STA A 0,X	ALL DONE
1009	EDU *	
1009 DE 3B	LDX SYMPTR	
100B B0 01 21	STX LOOKUP	
100E FF 10 22	JSR GNNLAS	
1011 09	DEX	
1012 EE 00	LDX	IS THIS REALLY NEW?
1014 9C 3B	CPX SYMPTR	
1016 2A A2	RNE GNM1	NOPE, START OVER
1018 FE 10 22	LDX GNNLAS	YES, RETURN WITH NEW ATOM
101B 39	RTS	
101C	GNMHAD EDU *	
101C CE 10 EA	LDX MENMENS	'NO MORE ROOM FOR TEMP NAME'
101F 7E 01 36	JMP ERREX	
1022 00 00	GNNLAS FDB 0	

NO ERROR(S) DETECTED

SYMBOL TABLE:

108D 39	RTS	
108E EDU	PTCDBR	REPORT ERROR
108E B0 71 3C	JSR DOSRPT	
1091	EDU *	
1091 B0 0E FB	JSR CLOSFI	CLOSE FILE
1094 20 F4	BRA PTCRT	AND RETURN
1096	EDU *	
109A 4C	FCC 'Load'	
109A 00	FDB 0	
109B 0E B3	FDB LOOKUP	
109B AF	FCC 'open'	
10A1 00	FDB 0	
10A2 0E 00	FDB OPENUM	
10A4 AF	FCC 'open'	
10A9 00	FDB 0	
10AA 0E F2	FDB OPENUM	
10AC 00	FDB 0	
10AD 43	EDU *	
10AD 43	FCC 'close'	
10B2 00	FDB 0	
10B3 0F 00	FDB CLSFUN	
10B5 43	FCC 'Close'	
10B8 00	FDB 0	
10BC 0E F9	FDB CLOFUN	
10BE 72	FCC 'return'	
10C3 00	FDB 0	
10C4 0F B3	FDB RATFUN	
10C4 67	FCC 'getname'	
10C8 00	FDB 0	
10CE 0F B8	FDB GNMFUN	
10D0 00	FDB 0	
10D1 42	FCC 'BAD ARG TO OPEN/LOAD'	
10E5 04	FDB 4	
10E4 4E	FCC 'NO SPACE FOR NEW BENAME ATOM'	
1103 04	FDB 4	
1104	EDU *	
1104 00 00	FDB 0	
1104 00 00	FDB 0	
1108	RMB 24	
1120 00 00	FDB 0	
1122	RMB 162	
11C1	BEGADR EDU *	

11E0 OPT OBEAD MAG SUBRS2 ENDS AT \$11C4
 0000 EQU 0
 0002 EQU 2
 0004 EQU 4
 * DIRECT PAGE STORAGE CELLS
 0020 ORG \$20
 0020-13 E5 BEGRD THIS SHOULD BE IN EACH MODULE
 0022 STR1 RMB 2
 0024 STR2 RMB 2
 0026 STP1 RMB 2
 0028 STP2 RMB 2
 002A XIMP2 EQU STP1
 002C XIMP2 EQU STP2
 002E FREPTR RMB 2
 0030 SFCPTR RMB 2
 0032 ENDFTR RMB 2
 0034 SYNLST RMB 2
 0036 MANPTR RMB 2
 0038 FURA RMB 2
 003A LSTPTR RMB 2
 003C SYMTR RMB 2
 003E CELPTR RMB 2
 0040 SFSAVE RMB 2
 0042 STUPTR RMB 2
 0044 CNPTM2 RMB 2
 0046 ALP RMB 2
 0048 HLP RMB 2
 004A HILATM RMB 2
 004C LPARAT RMB 2
 004E RPARAT RMB 2
 0050 DOTATM RMB 2
 0052 SQUATM RMB 2
 0054 QUATM RMB 2
 0056 LANATM RMB 2
 0058 ALANAT RMB 2
 005A NSURAT RMB 2
 005C SUBRAT RMB 2
 005E TATON RMB 2
 0060 COMATM RMB 2
 0062 ASETAT RMB 2
 0064 EDIATH RMB 2
 0066 SUBLS2 RMB 2
 0068 LBRKAT RMB 2
 006A RBRKAT RMB 2
 006C CUREVL RMB 2
 006E SURPTR RMB 2
 0070 BABPTR RMB 2
 0072 GCTEAP RMB 2
 0074 GCFREE RMB 2
 0076 FBR SUBLS3
 * SINGLE CHAR SLOTS ...
 00F0 ORG \$F0
 00F0 PEEKC RMB 1
 00F1 RSTFLO RMB 1
 * GLOBAL CONSTANTS
 00FC ENDRM RMB 2
 00FE * ZERO WORD
 * GLOBAL VECTORS
 0100 ORG \$100
 0100 EVAL RMB 3
 0103 READ RMB 3
 0106 PRINT RMB 3
 0109 GETCEL RMB 3
 010C

010F FRECEL RMB 3
 0112 PUSHX RMB 3
 0115 POPX RMB 3
 0118 TOPX RMB 3
 011B EVALM RMB 3
 011E SETATM RMB 3
 0121 LOOKUP RMB 3
 0124 GAXTL RMB 3
 0127 GFAXTL RMB 3
 012A LSTINT RMB 3
 012D LSTAD0 RMB 3
 0130 LSTAD2 RMB 3
 0133 LSTEND RMB 3
 0136 ERREX RMB 3
 0139 ATININT RMB 3
 013C ISDIPR RMB 3
 013F ISATON RMB 3
 0142 BETC RMB 3
 0145 JMP NUMINW
 0148 MURFM RMB 3
 014B FRNUM RMB 3
 014E PUTSYH RMB 3
 0151 PUTC RMB 3
 0154 ERBRK RMB 3
 0157 GETSYH RMB 3
 015A STKFRG RMB 3
 015D GCOL RMB 3
 0160 PROFSH RMB 3
 0163 GETCAR RMB 3
 0166 ISVAR RMB 3
 0169 PRIMR RMB 3
 016C PROPOP RMB 3
 *
 11E0 ORG OBEAD
 2100 MAGURD EQU \$2100 ('-')
 11E0 ATMFUN EQU *
 11E0-21 00 MAGURD
 * (ATOM X) RETURNS TRUE IF X EVALS TO ATOM, NUMBER, OR NIL
 11E2 8D 6F BSR GNXTA
 11E4 8D 01 3F JSR ISATON
 11E7 25 03 RCS FALSE
 11E9 EQU *
 11E9 DE SE TRUE
 11EB 39 RTS
 11EC DE FE FALSE
 11EE 39 RTS
 *
 11EF EQU *
 11EF 21 00 FDB MAGURD
 11F1 8D 60 BSR GNXTA
 11F3 8D 01 3F JSR ISATON
 11F6 25 F4 RCS FALSE
 11F8 2C F2 BGE FALSE
 11FA 20 ED BRA TRUE
 *
 11FC ADDFUN EQU *
 11FC 21 00 FDB MAGURD
 11FE 8D 31 BSR GNXTM
 1200 DF 48 STX FLP
 1202 8D 20 BSR GNXTM
 1204 EQU *
 1204 DF 4A STX MLP
 1206 9A 49 LDA A FLP+1
 1208 9B 47 ADD A MLP+1
 120A 19 DAA

NUMBER OF FREE CELLS AFTER LAST GCOL

INITIALIZED IN LISP

1208 97 49 STA A FLP+1
 1209 96 48 LDA A FLP
 120F 99 4A ADC A MLP
 1211 19 DAA
 1212 97 48 STA A FLP
 1214 0E 48 LDX FLP
 1216 7E 01 48 JMP NUMFRM
 1219 EQU *
 1219 21 00 SUBFUN EQU *
 121B SUBFNI EQU *
 121B 8D 14 BSR GNXTM
 121D DF 48 STX FLP
 121F 8D 10 BSR GNXTM
 1221 8D 12 E5 JSR NUMINW
 1224 20 0E BRA ADDF2
 *
 1226 GTRFUN EQU *
 1226 21 00 FDB MAGURD
 1228 8D F1 BSR SUBFNI
 122A 8D 01 4B JSR FRNUM
 122D 2F 80 BLE FALSE
 122F 20 88 BRA TRUE
 *
 1231 GNXTM EQU *
 1231 DE 44 LDX ALP
 1233 8D 01 24 JSR GNXTL
 1236 DF 44 STX ALP
 1238 DE 26 LDX XTMP
 123A EQU *
 123A 8D 01 3F JSR ISATON
 123D 25 05 RCS GNHRM
 123F 2C 03 BGE GNHRM
 1241 7E 01 48 JMP FRNUM
 1244 *
 1244 DF 34 STX FORN
 1246 CE 13 CB LDX GNHRMS
 1248 8D 01 54 JSR ERBRK
 124C 96 F1 LDA A RSTFLO
 124E 27 EA BEQ GMR2
 1250 DE FE LDX ZERO
 1252 39 RTS
 *
 1253 GNXTA EQU *
 1253 DE 44 LDX ALP
 1255 8D 01 24 JSR GNXTL
 1258 DF 44 STX ALP
 125A DE 26 LDX XTMP
 125C 39 RTS
 *
 125D CNMUN EQU *
 * (CONCURRENCE (....) (LIST ...)) (NOTHING))
 * REPEATEDLY EVALS FORN, CONCATS RESULTS TOGETHER.
 * STOPS WHEN RESULT IS FOTATM (I.E. NOTHING)
 125F 8D 01 2A JSR LSTINT
 1262 8D EF FDB MAGURD
 1264 25 1E BSR GNXTA
 1266 DF 48 BCS CNMUNL
 1268 EQU *
 1268 8D 01 00 JSR EVAL
 126A 8D 01 3C JSR ISDIPR
 126E 27 07 BEQ CNMUNL
 1270 25 09 BCS CNMUNL
 1272 8D 12 95 JSR LSTCON
 1275 25 0B RCS CNMUN2
 1277 EQU *
 1277 DE 48 LDX FLP
 1279 20 ED BRA CNMUNL AND GO RE-EVAL IT

1278 9C 64 CNMUN EQU *
 1278 CPX EDIATH
 127B 27 05 BEQ CNMUN2
 127E 8D 01 33 JSR LSTEND
 1282 20 03 BRA CNMUN3
 1284 EQU *
 1284 8D 01 15 JSR LSTENO
 1287 EQU *
 1287 7E 01 6C CNMUN3 JMP PKOPOP
 *
 128A LSTFUN EQU *
 128A → (LIST 'A' 'B' 'C')
 128A 21 00 FDB MAGURD
 128C DE 3E LDX STKPTR
 128E 4F 00 CLR CAR.X
 1290 4F 01 CLR CAR+1,X
 1294 39 RTS
 *
 1295 LSTCON EQU *
 * CONCAT A LIST TO LIST AT TOP OF STACK
 * RETURN WITH C-BIT SET IF CDR OF LAST DTPR NON-NIL
 1295 DF 34 STX FORN
 1297 8D 01 18 JSR TOPX
 129A 96 34 LDA A FORN
 129C A7 02 STA A CDR.X
 129E 9A 35 LDA A FORM+1
 12A0 A7 03 STA A CDR+1,X
 12A2 EQU *
 12A2 DF 28 STX XTMP2
 12A4 EE 02 LDX CDR.X
 12A6 8D 01 3C JSR ISDIPR
 12A9 24 F7 BCC LSTC2
 12AB 28 01 BNE LSTC3
 12AC 0C CLC
 12AE UE 3E LDX STKPTR
 12B0 96 28 LDA A XTMP2
 12B2 A7 00 STA A CAR.X
 12B4 96 29 LDA A XTMP2+1
 12B6 A7 01 STA A CAR+1,X
 12B8 DE 34 LDX FORN
 12BA 39 RTS
 *
 12B9 EWLSFN EQU *
 * (EVALLIST (CDR L))
 * EVAL ALL ELEMENTS OF ARG (ASSUMED TO BE A LIST)
 * RETURN WITH VALUE OF LAST ONE
 12B9 DE 44 LDX ALP
 12BB 8D 01 3C JSR ISDIPR
 12C2 25 08 BCS EVLLS2
 12C4 EE 00 LDX CAR.X
 12C6 20 04 BRA EVLLS2
 *
 12C8 PGMFN EQU *
 * (PRGM A B C D)
 * EVALS ALL ARGS AND RETURNS THE VALUE OF THE LAST ONE
 12C8 21 00 LDX FDB
 12CA DE 44 LDX ALP
 12CC EQU *
 12CC DF 48 JSR FLP
 12CE EQU *
 12CE 8D 01 3C JSR ISDIPR
 12D1 25 0F BCS PGMFN
 12D3 DF 44 LDX ALP
 12D5 EE 00 LDX CAR.X
 12D7 8D 01 00 JSR EVAL

CLEAR TOP OF STACK (TO AVOID FREEING LIST)

RETURN WITH ARG LIST

SAVE POINTER TO LIST

FILL IN NEW CDR

SAVE POINTER TO LAST CELL

NOU POINT TO NEXT CELL IN LIST

IS IT A DTPR?

YES, KEEP LOOKING

SKIP THE CLL IF NON-NIL CDR

FILL IN NEW VALUE FOR LIST POINTER

RESTORE X-NEG

RETURN WITH C-BIT CORRECT

NOT A DTPR, JUST RETURN THIS

GET ARG (SHOULD BE A LIST)

AND CONTINUE WITH PRGM CODE

INITIALIZE RETURN VALUE

EVAL CAR

128A BF 4B STX FLP
 128C DE 44 LDX ALP
 128E EE 02 LDX CRD.X
 12E0 20 EC BRA PGMFN2
 12E2 EDU *
 12E2 DE 4B LDX FLP
 12E4 3F RTS

12E5 NUMINV EDU *
 * EXPECTS NUMBER IN X-REG (BCD 10'S COMPLEMENT)
 * PRODUCES 10'S COMPLEMENT

12E5 BF 24 STX XTHP
 12E7 27 14 BEO MNIVRT
 12E9 86 99 LDA A #999
 12ER 16 TAB
 12EC 80 26 SUB B XTHP
 12EE 90 27 SUR A XTHP+1
 12F0 8B 01 ARD A B1
 12F2 19 BAA
 12F3 97 27 STA A XTHP+1
 12F5 17 TRA
 12F6 89 00 ADC A #0
 12F8 19 BAA
 12F9 97 26 STA A XTHP
 12FB DE 26 LDX XTHP
 12FD MNIVRT EDU *
 12FD 3F RTS

12FE RSTFUN EDU *
 12FE 21 00 FDB MAGURD
 1300 RSTFN2 EDU *
 1300 86 FF LDA A B-1
 1302 RSTFN3 EDU *
 1302 97 F1 STA A RSTFLG
 1304 RTNEDI EDU *
 1304 DE 64 LDX EDIATH
 1306 RTFRM EDU *
 1306 3F RTS

1307 RTBFUN EDU *
 1307 21 00 FDB MAGURD
 1309 8B 12 53 JSR GNXTA
 130C 8D 01 3F JSR ISATOM
 130F 25 0B BCS RTBFN2
 1311 2C 08 RGE RTBFN2
 1313 8D 01 4B JSR FRHNUM
 1316 2F EB BLE RSTFN2
 1318 BF 26 STX XTHP
 131A 76 27 LDA A XTHP+1
 131C 20 E4 BRA RSTFN3

131E RTBFN2 EDU *
 131E 86 CF LDA A #FFF-'0
 1320 20 E0 BRA RSTFN3

1322 BRKFUN EDU *
 1322 21 00 FDB MAGURD
 1324 DE FE LDX ZERO
 1326 DF 34 STX FORN
 132B CE 13 DB LDX #BRKNSG
 132B 7E 01 54 JMP ERRBRK

132E RTBFUN EDU *
 132E 21 00 FDB MAGURD
 1330 7F 00 49 CLR
 1333 8D 12 53 JSR GNXTA
 1336 8D 01 3F JSR ISATOM
 1339 25 07 BCS RTBFUN

133R 2C 05 RTFUN1
 133D 8D 01 4R JSR FRHNUM
 1340 DF 4B STX FLP
 1342 BITFUN1 *
 1342 DE AC LDX CUREVL
 1344 DF 44 STX ALP
 1346 BITFUN2 *
 1346 8D 12 53 JSR GNXTA
 1349 25 89 BCS RTNEDI
 134B 96 49 LDA A FLP+1
 134D 8B 99 ADD A #999
 134F 19 BAA
 1350 97 49 STA A FLP+1
 1352 27 B2 REO RTFRM
 1354 8D 01 63 JSR GETCAR
 1357 8D 01 06 JSR PRINT
 135A 20 EA BRA BITFUN2

135C SUNLS3 *
 135C EDU *
 135C 61 FCC 'alon'
 1360 00 FCA 0
 1361 11 E0 FDB ATHFUN
 1363 8E FCC 'number'
 1369 00 FCC 0
 136A 11 EF FDB NUMFUN
 136F 00 FCC 'adj'
 1370 11 FC FDB 0
 1372 73 FCC 'sub'
 1375 00 FCC 0
 1376 12 19 FDB SURFUN
 137B 47 FCC 'greater'
 137F 00 FCC 0
 1380 12 24 FDB GTRFUN
 1382 6C FCC 'list'
 1386 00 FCC 0
 1387 12 8A FDB LSTFUN
 1389 65 FCC 'evallist'
 1391 00 FDB 0
 1392 12 8B FDB EVLSFN
 1394 72 FCC 'rebrk'
 139A 00 FDB 0
 139B 13 07 FDB RTBFUN
 139B 62 FCC 'bl'
 139F 00 FDB 0
 13A0 13 2E FDB BITFUN

13A2 00 FDB 0
 13A3 MSURL3 *
 13A3 70 FCC 'progn'
 13A8 00 FDB 0
 13A9 12 C8 FDB PGMFUN
 13AB 43 FCC 'concatn'ile'
 13B4 00 FDB 0
 13B5 12 5B FDB CNUFUN
 13B7 72 FCC 'reset'
 13B8 00 FDB 0
 13B9 12 FE FDB RSTFUN
 13BF 62 FCC 'break'
 13C4 00 FDB 0
 13C5 13 22 FDB BRKFUN

13C7 00 FDB 0
 13C8 4E CNMENS FCC 'NUMBER REQUIRED'
 13D7 04 FDB 4
 13B8 42 BRKNSG FCC 'BREAKING ...'
 13E4 04 FDB 4
 13E5 BEGARU EDU *
 BEGARU END

133R 2C 05 PRINT OUT CAR'S OF CUREVL LIST
 RETURN EOI 50 NO PRINT OUT
 SUBTRACT I (BCD-U)ISE

IF ARG IS A NUMBER, RETURN FUN THAT FAR BACK ON CUREVL

NO ERROR(S) DETECTED

SYMBOL TABLE:

ADDU2 1204 ADDFUN 11FC ALP 0044 ATHFUN 11E0 ATHINI 0139
 BEGARU 13E5 BEGTR 0020 BRKFUN 1322 BRKNSG 13B8 BIFRIN 1306
 BITFUN 132E BITFUN2 1342 CAR 0000 CAR 0002
 CELPTR 003A CAPTR2 0042 CAPTRP 0040 CNUDW2 1284 CNUDW3 1287
 CNUDW 127B CUREVL 1268 CUFUN 1250 CUNXT 1277 COMATN 0040
 CUREVL 006C DADPTR 0070 DOTATA 0050 ENDHEN 00FC ENDPTR 002E
 EDI 0004 EDIATH 0064 ERBRK 0154 ERREX 0136 EVAL 0100
 EVLATH 011B EVLS2 12CC EVLSFN 128B FALSE 11EC FLP 004B
 FURN 003A FRECEL 010F FREPTR 002A FRHNUM 014B GCFREE 0074
 GCOL 015D GCTEMP 0072 GETC 0142 GETCAR 0163 GETCEL 010C
 GETSYN 0157 GENXTL 0127 GHW2 123A GHWERS 13C8 GHWERR 1244
 GNATA 1253 GNATL 0124 GNXTM 1231 GTRFUN 1226 ISATOM 013F
 ISDIFR 013C ISVAR 0166 LAMATH 0056 LBRKAT 008B LOOKUP 0121

LPAKAT 004C LSTA02 0130 LSTADU 0120 LSTC2 12A2 LSTC1 12AF
 LSTCOM 1295 LSTENO 0115 LSTENR 0133 LSTFUN 128A LSTINI 012A
 LSTFR 0036 MAGURD 2100 MANPTR 0032 MATH 004A MLARAT 005B
 NLP 0046 NMLVKT 12FD NSUBAT 005C NSUBLJ 13A3 NUMFRN 014B
 NUMFUN 11EF NUMJMV 12E5 OBEGAD 11E0 PEKC 00FO PGMUDN 12E2
 PGMFN2 12CE PGMFUN 12CB POPX 0115 PRINR 0169 PRINT 0106
 PRGFUP 014C PROF'SH 0160 PUSHX 0112 PUTC 0151 PUTSYN 014E
 QUARTR 0054 RBRKAT 006A READ 0103 RPARAT 004E RSETAT 0062
 RSTFLG 00F1 RSTFN2 1300 RSTFN3 1302 RSTFUN 12FE RTBFN2 131E
 RTBFUN 1307 RTNEDI 1304 SETATH 011E SDMPTR 006E SPCPTR 002C
 SFSAVE 003C SOUTAT 0052 STKFRG 015A STKPTR 003E STP1 0026
 STP2 002B STP1 0022 STR2 0024 SUBFM1 121B SUBFUN 1219
 SUBLS2 0066 SUBLS3 135C SUBRAT 005A SYNLS1 0030 SYNPTR 003B
 TADW 005E TOPX 011B TRUE 11EY XTMP 0026 XTMP2 002B
 ZERO 00FE

